

Lustre System Administration Tutorial

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Outline

1:00PM - Configuration/Tuning: (Dustin - 30 min)

- General benchmarking
- Router/client/server tuning

1:30PM - Monitoring/Metrics: (Dustin - 1 hour)

- Performance monitoring
- Health monitoring

2:30PM - 15 minute break

2:45PM - When Things Go Wrong: (Rick - 1 hour)

- Lustre recovery
- Gathering debug data: kdump, lctl dk
- Network debugging
- Repairing filesystem issues

3:45PM - Other Useful Admin Info: (Rick - 30 min)

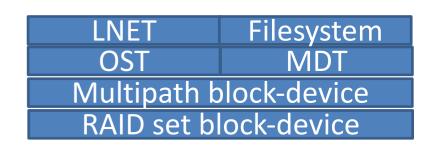
- Striping
- How RR and QOS allocator work
- Layouts
- PFL
- DoM

4:15PM – 45 minute panel session



Benchmarking

- Keep vendors accountable for requirements and system acceptance
- Establishing a performance baseline helps to detect performance issues later on.
- Benchmark at the different system layers:
 - Block-device/multipath-device (XDD/sgp_dd)
 - ZFS dRAID exception
 - LDISKFS/ZFS (fio)
 - LNET (Inet-selftest)
 - Filesystem data (obdfilter-survey)
 - Filesystem metadata (mdtest)
 - Filesystem data and metadata (IOR)



• There are many tools that can be used, these are just examples



Benchmarking block/multipath device

- Tuning can be different for block and multipath devices. Need to check both.
- Testing /dev/sd vs. /dev/sg can tell you how caching is impacting performance.
 - Need to understand connectivity to RAID system if IB/SAS/FC attached so you don't stress out the same LUN more than once
- XDD or sgp_dd are good tools to use
- XDD example:
 - /opt/xdd/bin/xdd.linux -rwratio 100 -targets 10 /dev/sg{0,2,4,6,8,30,32,34,36,38} -sgio -reqsize 1 -numreqs 10000000 blocksize 1048576 -verbose -passes 3 -queuedepth 16 -timelimit 120 seek truesequential -datapattern random



Benchmarking LDISKFS/ZFS

- You can use a tool like fio to write data to ldiskfs or ZFS to test your performance tunables
- Run streaming and random reads/writes
- Test metadata performance
 - Can impact the OST metadata lookup



Benchmarking LNET

- Lustre has a built in tool called Inet-selftest that will exercise the network only (no disk activity)
 - Pick which servers and clients to use
 - Reads/writes
 - Data size
 - Concurrency (simultaneous number of threads participating)
- This will allow you to ensure that the network is performing and tuned
- Very well documented here:
 - http://wiki.lustre.org/LNET_Selftest
 - We don't run this frequently, but we reference this documentation when we do (you don't have to memorize all this stuff)
 - Rick will be giving an example later



Benchmarking filesystem data layer

- Lustre has a built in tool called obdfilter-survey that will exercise the disk layer only (no LNET activity)
- This will allow you to ensure the disks are performing and tuned independent of LNET
- Note: This test can be destructive; run these tests before production
- More info on Lustre Wiki: http://wiki.lustre.org/OBDFilter_Survey



Benchmarking filesystem data layer

- Example:
 - modprobe obdecho
 - mkdir -p /tmp/obdfilter-survey_output
 - nobjlo=1 nobjhi=16 thrlo=1 thrhi=1024 size=32768 rslt_loc=/tmp/obdfilter-survey_output targets="testfs-OST0000 testfs-OST0001 testfs-OST0002" case=disk obdfilter-survey
- Tunables:
 - nobj[lo|hi]: concurrent object count per OST to iterate over
 - thr[lo|hi]: Thread range to iterate over
 - size: total amount of data to be written
 - Should target something that is >2x larger than WBC size of RAID controller
 - rslt_loc: Where to write the result data to
- You should write a script to coordinate the obdfilter-survey processes per node if you have a shared RAID sub-system (DDN, NetApp E-series, etc...)



Filesystem-level benchmarking

Node count:

- "Hero" performance (max performance possible)
 - Usually ~10-30% of total node count at large-scale
 - Node placement may be important (network layout, router layout, etc...)
- Single-client
 - Helps to understand what a small job will see
 - Helps to understand scaling behavior
- All clients (max performance from all clients)
 - Generally ~10-30% slower than "hero" at large-scale
 - Important to understand for full-scale jobs

Thread count:

- Single-threaded
 - Most common use case for users
- Multi-threaded
 - Max possible performance per node



Lustre filesystem metadata benchmarking

- mdtest is a very common tool for exercising filesystem metadata from multiple clients using MPI
- Considerations:
 - shared directory
 - unique directory
 - just metadata (zero-length files) or file IO too?

• Lots of tunables:

#-b: branching factor of hierarchical directory structure
#-B: no barriers between phases (create/stat/remove)
#-c: collective creates: task 0 does all creates and deletes
#-C: only create files/dirs
#-d: the directory in which the tests will run
#-D: perform test on directories only (no files)
#-e: number of bytes to read from each file
#-E: only read files
#-f: first number of tasks on which the test will run
#-F: perform test on files only (no directories)
#-h: prints help message
#-i: number of iterations the test will run
#-l: last number of tasks on which the test will run
#-L: files/dirs created only at leaf level

#-n: every task will create/stat/remove # files/dirs per tree
#-N: stride # between neighbor tasks for file/dir stat (local=0)
#-p: pre-iteration delay (in seconds)
#-r: only remove files/dirs
#-R: randomly stat files/dirs (optional seed can be provided)
#-s: stride between the number of tasks for each test
#-S: shared file access (file only, no directories)
#-t: time unique working directory overhead
#-T: only stat files/dirs
#-u: unique working directory for each task
#-v: verbosity (each instance of option increments by one)
#-V: verbosity value
#-w: number of bytes to write to each file
#-y: sync file after write completion
#-z: depth of hierarchical directory structure



Lustre filesystem metadata benchmarking

File creates, metadata only, unique-directory, 3 iterations, 5-minute delay

```
#BSUB -q storage# Job queue#BSUB -o mdtest_unique_dir_multi-node.o%J# output is sent to file job.output#BSUB -e mdtest_unique_dir_multi-node.e%J# error is sent to file job.error#BSUB -J mdtest_unique_dir_multi-node# name of the job#BSUB -nnodes 630# Number of nodes to use in the job#BSUB -W 360# wallclock -W [hour:]minute[/host_name | /host_model]#BSUB -U PT#BSUB -P ACCEPTANCE
```

```
MOUNT="alpine"
```

```
BINDIR="/gpfs/alpine/stf002/scratch/leverman/alpine_acceptance"

OUTDIR="$BINDIR/${LSB_JOBID}_md_test"

[ -e $OUTDIR ] || {

mkdir -p $OUTDIR

}

cd $BINDIR

module load gcc

jsrun -n 630 -c ALL_CPUS -a 20 -X 1 $BINDIR/build/mdtest_build/mdtest -n 32768 -p 300 -F -u -C -r -i 3 -v -v -u

$OUTDIR
```



Lustre filesystem data benchmarking

- IOR is a very common tool for exercising filesystem from multiple clients using MPI
- Considerations:
 - FPP or SSF
 - Random vs. Sequential workload (random is more realistic on an aging system)
 - picking IO size (alignment with RAID engine or user workload)
 - picking the amount of data (want to write for long enough to exceed client, server, and RAID engine caches
 - Don't let vendors stonewall, pre-create, etc... as part of acceptance

Lots of tunables:

-a S api -- API for I/O [POSIX|MPIIO|HDF5|HDF5|S3|S3_EMC|NCMPI] -A N refNum -- user reference number to include in long summary

- -b N blockSize -- contiguous bytes to write per task (e.g.: 8, 4k, 2m, 1g) -B useO DIRECT -- uses O DIRECT for POSIX, bypassing I/O buffers
- -c collective -- collective I/O
- -C reorderTasksConstant -- changes task ordering to n+1 ordering for readback
- -d N interTestDelay -- delay between reps in seconds
- -D N deadlineForStonewalling -- seconds before stopping write or read phase
- -e fsync -- perform fsync upon POSIX write close
- -E useExistingTestFile -- do not remove test file before write access
- -f S scriptFile -- test script name
- -F filePerProc -- file-per-process
- -g intraTestBarriers -- use barriers between open, write/read, and close
- -G N setTimeStampSignature -- set value for time stamp signature
- -h showHelp -- displays options and help
- -H showHints -- show hints
- -i N repetitions -- number of repetitions of test
- -I individualDataSets -- datasets not shared by all procs [not working]
- -j N outlierThreshold -- warn on outlier N seconds from mean
- -J N setAlignment -- HDF5 alignment in bytes (e.g.: 8, 4k, 2m, 1g)
- -k keepFile -- don't remove the test file(s) on program exit
- -K keepFileWithError -- keep error-filled file(s) after data-checking
- -I data packet type-- type of packet that will be created [offset|incompressible|timestamp|o|i|t]
- -m multiFile -- use number of reps (-i) for multiple file count

-M N memoryPerNode -- hog memory on the node (e.g.: 2g, 75%)

- -n noFill -- no fill in HDF5 file creation
- -N N numTasks -- number of tasks that should participate in the test
- -o S testFile -- full name for test
- -O S string of IOR directives (e.g. -O checkRead=1,lustreStripeCount=32)
- -p preallocate -- preallocate file size
- -P useSharedFilePointer -- use shared file pointer [not working]
- -q quitOnError -- during file error-checking, abort on error
- -Q N taskPerNodeOffset for read tests use with -C & -Z options (-C constant N, -Z at least N) [!HDF5]
- -r readFile -- read existing file
- -R checkRead -- check read after read
- -s N segmentCount -- number of segments
- -S useStridedDatatype -- put strided access into datatype [not working]
- -t N transferSize -- size of transfer in bytes (e.g.: 8, 4k, 2m, 1g)
- -T N maxTimeDuration -- max time in minutes to run tests
- -u uniqueDir -- use unique directory name for each file-per-process
- -U S hintsFileName -- full name for hints file
- -v verbose -- output information (repeating flag increases level)
- -V useFileView -- use MPI_File_set_view
- -w writeFile -- write file
- -W checkWrite -- check read after write
- -x singleXferAttempt -- do not retry transfer if incomplete
- -X N reorderTasksRandomSeed -- random seed for -Z option
- -Y fsyncPerWrite -- perform fsync after each POSIX write
- -z randomOffset -- access is to random, not sequential, offsets within a file
- -Z reorderTasksRandom -- changes task ordering to random ordering for readback



Lustre filesystem data benchmarking

FPP, read/write, I6MB transfer size

#!/bin/bash
#BSUB -q storage # Job queue
#BSUB -o IOR_fpp_32MB_seq_alpine.o%J # output is sent to file job.output
#BSUB -e IOR_fpp_32MB_seq_alpine.e%J # error is sent to file job.error
#BSUB -J IOR_fpp_32MB_seq_alpine # name of the job
#BSUB -nnodes 504 # Number of nodes to use in the job
#BSUB -W 240 # wallclock -W [hour:]minute[/host_name | /host_model]
#BSUB -P ACCEPT
#BSUB -alloc_flags "smt4 isolategpfs" # Isolate GPFS processes and configure for SMT4

MOUNT=\$(pwd | awk -F/ '{print \$3}') BDIR="/gpfs/alpine/stf002/scratch/leverman/alpine_acceptance" TDIR="\$BDIR/ior_testdir" ITERS=3 BSIZE="7168g" INTERFACE="POSIX" TSIZE="16m"

mkdir -p \${TDIR} cd \${BDIR} module load gcc

date

exit 0

echo "POSIX read/write run for seq file per process 16MB transfer size, 20min" jsrun -n 504 -c ALL_CPUS -a 1 \${BDIR}/build/ior-3.1.0/src/ior -g -d 360 -o \${TDIR}/POSIX_fpp_ior -F -i \${ITERS} -b \${BSIZE} -t \${TSIZE} -w -r -a \${INTERFACE} -e -v -v date



Lustre Tuning (general – ALL)

- Tuning BIOS
 - Disable c-states
 - Put in "performance" mode
 - Performance power governor in OS
- Ko2iblnd (our lustre file systems are all IB attached that will be the assumption for these slides)
 - options ko2iblnd ib_mtu=2048 timeout=100 credits=2560 ntx=5120 peer_credits=63 concurrent_sends=63 fmr_pool_size=1280 fmr_flush_trigger=1024
- LNET
 - /etc/modprobe.d/Inet.conf
 - options lnet check_routers_before_use=1 router_ping_timeout=120 dead_router_check_interval=50 avoid_asym_router_failure=0 live_router_check_interval=50
 - /etc/Inet.conf
 - net:
 - net type: o2ib2
 - local NI(s):
 - - nid:
 - interfaces:
 - 0: ib1
 - tunables:
 - peer_timeout: 180
 - peer_credits: 63
 - peer_buffer_credits: 0
 - credits: 2560
 - global:
 - discovery: 0
 - Peer Credits
 - Modern systems you generally set peer credits to 63 (may need to be lower with FDR IB 8 because of concurrent sends issue)
 - Compute vendors may set something specific (need to keep credit the same across all clients and servers)
- Striping (talk about this later)





Lustre Server Tuning

- SRP (for an SFA14KX):
 - /etc/modprobe.d/ib_srp.conf
 - options ib_srp cmd_sg_entries=255 indirect_sg_entries=2048 allow_ext_sg=1 use_blk_mq=N
- Block devices (udev rules for an SFA14KX):
 - KERNEL=="sd*", ENV{ID_VENDOR}=="DDN*", ENV{ID_MODEL}=="SFA14KX*", \
 - ATTR{device/timeout}="68", \
 - ATTR{queue/scheduler}="deadline", \
 - ATTR{queue/nr_requests}="192", \
 - ATTR{queue/read_ahead_kb}="0", \
 - ATTR{queue/max_sectors_kb}="\$attr{queue/max_hw_sectors_kb}"
 - KERNEL=="dm-*", ACTION=="change", ENV{NCCS_DM_TABLE}=="multipath" \}
 - ATTR{queue/scheduler}="deadline",
 - ATTR{queue/nr_requests="192}",
 - ATTR{queue/read_ahead_kb}="0}",
 - ATTR{queue/max_sectors_kb}="8192"
- Multipathd (for an SFA14KX):
 - device {
 - vendor "DDN"
 - product "SFA14KX"
 - prio "alua"
 - prio_args "exclusive_pref_bit"
 - path_grouping_policy "group_by_prio"
 - path_checker "tur"
 - path_selector "round-robin 0"
 - rr_weight "uniform"
 - failback "2"
 - no_path_retry "12"
 - user_friendly_names "yes"
 - dev_loss_tmo "10"
 - fast_io_fail_tmo "5"
 - max_sectors_kb "8192"
 - •
- LDISKFS/ZFS tunables:
 - options zfs metaslab_debug_unload=1 zfs_arc_max=15000000000 zfs_vdev_scheduler=deadline zfs_prefetch_disable=1 zfs_dirty_data_max_percent=30 zfs_dirty_data_max_max=60236916326 zfs_dirty_data_max=60236916326 zfs_arc_average_blocksize=2097152 zfs_max_recordsize=2097152 zfs_vdev_aggregation_limit=2097152 zfs_multihost_interval=60000



LDISKFS/ZFS Multipath block-device RAID set block-device SRP/iSER/SAS RAID subsystem

Lustre Client Tuning

- lctl set_param osc.*.checksums=0
 - You may already have network checksums enabled and don't need this
 - Performance penalty
- lctl set_param timeout=600
- lctl set_param ldlm_timeout=200
- lctl set_param at_min=250
- lctl set_param at_max=600
- lctl set_param ldlm.namespaces.*.lru_size=128
 - Might be ignored sometimes (current bug LU11518)
 - Low number for computes (generally), high number for login nodes
- lctl set_param osc.*.max_rpcs_in_flight=32
- lctl set_param osc.*.max_dirty_mb=64
- lctl set_param debug="+neterror"
 - Rick will talk more about this later



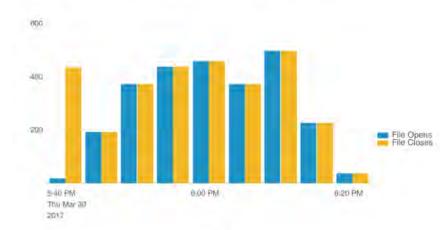
Lustre Router Tuning

- Check if LNET routing is enabled on this node
 - cat /sys/kernel/debug/lnet/routes
- LNET router buffer sizes
 - Defaults are generally too small
 - Can be changed on the fly
 - How we tune it:
 - tiny: 8192
 - Zero-payload (signals and acks)
 - small: 131072
 - 4k payload (metadata, zero-length file, etc...)
 - large: 4096
 - 1m max payload (file data)
- Different credit sizes per interface?
 - Depends on the networks you are routing between

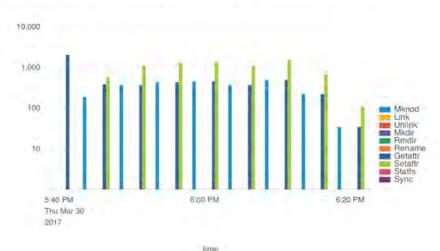


- Jobstats (job-level)
 - We assume single job per node and tag each lustre client with a job ID using the scheduler prologue/epilogue
 - You can gather this data as time-series or just have a report for what the total IO activity for the job was using tools like splunk/influx-Grafana

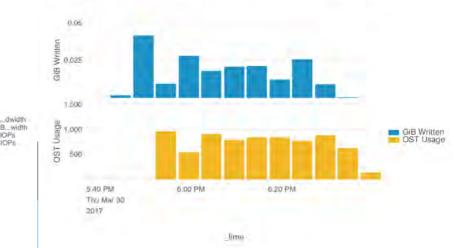
Job Specific I/C	Statistics: File	e Opens &	Closes
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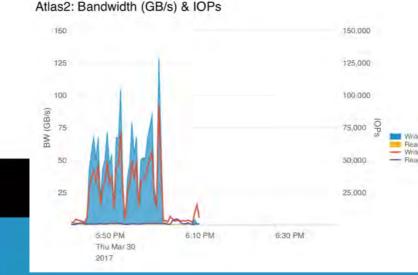
Job Specific I/O Statistics: Other Metadata Operations



Job Specific I/O Statistics: Write BW & OST Usage



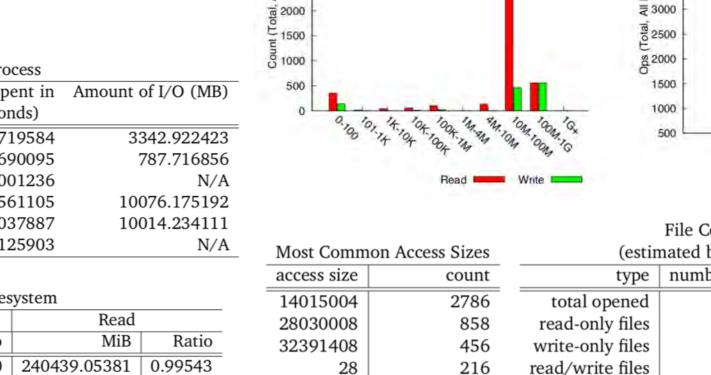
- job_id:	Titan-titan-login1							
<pre>snapshot_time:</pre>	1556021522							
read_bytes:	{ samples:	2,	unit:	bytes, min:	32768, max:	1015808,	sum:	1048576 }
write_bytes:	{ samples:	0,	unit:	bytes, min:	0, max:	0,	sum:	0 }
getattr:	{ samples:	0,	unit:	reqs }				
setattr:	{ samples:	0,	unit:	reqs }				
punch:	{ samples:	0,	unit:	reqs }				
sync:	{ samples:	0,	unit:	reqs }				
destroy:	{ samples:	0,	unit:	regs }				
create:	{ samples:	0,	unit:	reqs }				
statfs:	{ samples:	0,	unit:	reqs }				
get_info:	{ samples:	0,	unit:	reqs }				
set_info:	{ samples:	0,	unit:	reqs }				
quotact1:	{ samples:	0,	unit:	regs }				
[root@atlas-oss1a	1 ~]#							





- Darshan (job-level)
 - Load an environment module
 - Users compile code with darshan loaded
 - Darshan intercepts I/Os and gathers statistics
 - Tools exist to visualize data
 - Minimal performance impact

	Average I/O per process	
	Cumulative time spent in	Amount of I/O (MB)
	I/O functions (seconds)	
Independent reads	10.719584	3342.922423
Independent writes	1.690095	787.716856
Independent metadata	0.001236	N/A
Shared reads	28.561105	10076.175192
Shared writes	46.037887	10014.234111
Shared metadata	0.125903	N/A



Average I/O cost per process

POSIT

I/O Sizes

Mp.

100

80

60

40

20

0

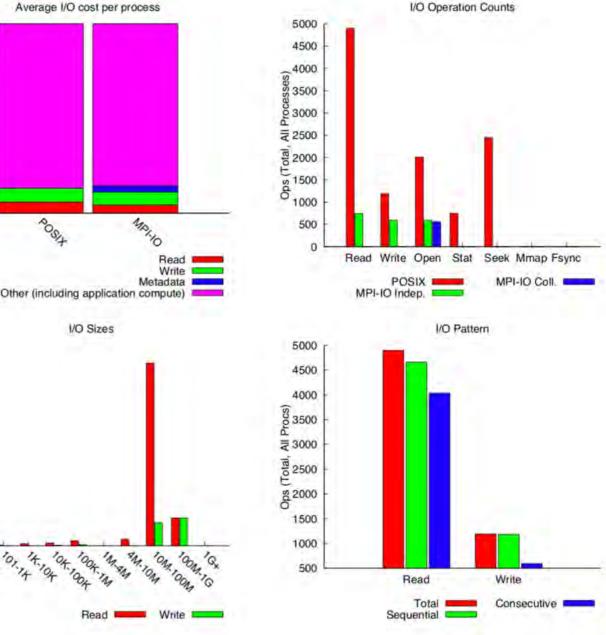
4000

3500

8 3000

ã 2500 T

Percentage of run time



		File Count Summ	nary	
zes	(estin	nated by I/O acce	ss offsets)	
ınt	type	number of files	avg. size	max size
86	total opened	18	13G	177G
58	read-only files	5	8.0G	20G
56	write-only files	6	2.4G	4.7G
16	read/write files	1	177G	177G
	created files	7	28G	177G

	Data Transf	er Per Files	system	
File System	Write		Read	
File System	MiB	Ratio	MiB	Ratio
/lustre/atlas1	194435.11740	1.00000	240439.05381	0.99543
/lustre/atlas	0.00000	0.00000	1104.70326	0.00457

- Brw_stats (OST/MDT target level)
 - Can use a data collector tool (cerebrod, telegraf, etc...) to collect the brw_stats data for each OST
 - Put this data into analytics tool (like splunk) to visualize
 - Dumps the following data:
 - pages per bulk r/w
 - discontiguous pages
 - disk I/Os in flight
 - I/O time (1/1000s)
 - disk I/O size

									-	54.
Q New Search								Save A	s 🗸 Close	64:
index=brw_stats "atlas*" datatype=:	io_size (rc_delta > 0 (<pre>MR wc_delta > 0) eval</pre>	bucket = if(bucket="	M", "1024K", bucket)	chart sum(rc_delta) a	s "Read" sum(wc_delta) as	"Write" over bucket	Year to da	te v Q	
64,101,048 of 80,236,585 events matched No	Event Sampling 🛩						Job v II		mart Mode 🛩	128:
Events Patterns Statistics (9)	Visualization								2	256:
al Column Chart / Format III Trellis										512:
4,000,000,000										
										1K;
3,000,000,000										2K:
2,000,000,000									4	4K:
									Read	
1,000,000,000									1000	8K:
										16K:
			-		10.04					LOK.
4K	SK	16K	32K	64K bucket	128K	256K	512K	1024K		32K:
T. A. A.					Read					
bucket					1072855875				994145671	64K:
Bei					145257286				and the second se	12017
104					51987006				56057424	128K:
928					45372995				74877889	256K:
CAN .					37937329				59397702	200N.
128K1					53114467				47675277	512K:
256K					253381278					JILIN.
512K					200124529				36139551	1M:
10248					2859735421				3786261580	

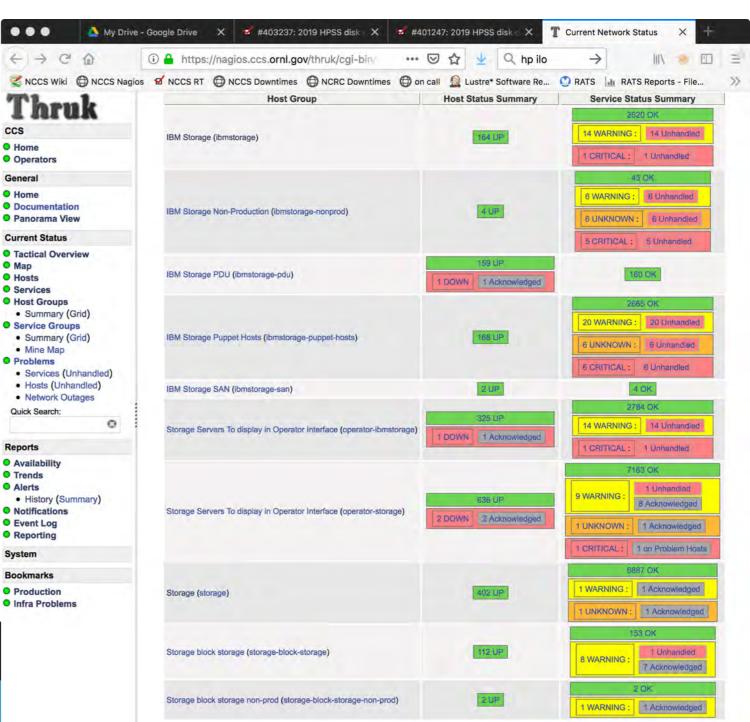
		read	d	1	wr	ite	
disk I/O size	ios	%	cum	%	ios		% cum %
1:	0	0	0	1	13087	0	0
2:	0	0	0	1	197	0	0
4:	0	0	0	- 1	573	0	0
8:	7	0	0	- 1	391	0	0
16:	0	0	0	1	13275	0	0
32:	1	0	0	1	7211	0	0
64:	1	0	0	- 1	12145	0	0
128:	4249	0	0	1	35584	0	0
256:	1	0	0	- 1	29320	0	0
512:	1	0	0	1	133477	0	0
1K;	0	0	0	- 1	446480	0	0
2K:	0	0	0	- 1	984317	0	1
4K:	105868354	35	35	- 1	391162	3	3 4
8K:	1964299	0	36	- 1	129412	02	10 15
16K:	8549687	2	39	1	485714	3	4 19
32K:	8640700	2	42	- 1	349431	4	2 22
64K:	10565910	3	45	1	107824	9	0 23
128K:	4533522	1	47	1	163948	1	1 24
256K:	4850565	1	49	1	183808	3	1 26
512K:	37528509	12	61	1	130405	9	1 27
1M:	113179958	38	100		879366	19	72 100

- LMT Itop
 - Collects metadata, bandwidth, and other server-side stats
 - Puts data in a database via data collection tool (cerebrod)
 - Different interfaces to view the data (Itop and Iwatch)
- Controller-local IO statistics
 - DDN, NetApp, Adaptec, etc... should present B/W, I/O size, IOPS, latency etc... for LUNs, PDs, host ports, etc...

Files	system:	atlasi											
	Inodes:			total		467.689	n used	(46%),	5	56.3	l1m fi	ree	
	Space:							(73%),			97t fi		
By	tes/s:	1	L.654g	read,		1.533	g write	e,		66	598 I	OPS	
ME)ops/s:	29188	B open,	2	6767	close,	975	1 getatt	r,	1356	6 set	attr	
			link,			unlink,		4 mkdir,					
						rename,		0 getxat					
>0ST	Ś	055						LOCKS		LCR	%сри	%mem	%spc
(7)	las-o	oss1a1	20058	0	11	1	45	57594	28	21	6		73
(7)	las-c	oss1a2	20058	0	8	2	35	54866	162	24	6	98	74
(7)	las-c	oss1a3	20058	0	8	2	31	53153	68	58	5	99	75
(7)	las-c	oss1a4	20058	0	6	2	36	57998	42	34	6	99	73
(7)	las-a	oss1a5	20058	0	60	0	79	55888	38	28	6	99	74
(7)	las-a	oss1a6	20058	0	6	4	39	55211	43	30	6	100	73
(7)	las-c	oss1a7	20058	0	51	2	69	56434	38	26	6	100	73
(7)	las-c	oss1a8	20058	0	6	1	42	56182	77	55	6	99	73
(7)	las-c	oss1b1	20058	0	7	2	29	60166	42	29	6	99	73
(7)	las-c	oss1b2	20058	0	49	2	82	55486	57	41	6	100	74
(7)	las-c	oss1b3	20058	0	6	31	64	56988	37	27	6	100	73
(7)	las-c	oss1b4	20058	0	7	1	27	54338	39	29	6	99	75
(7)	las-c	oss1b5	20058	0	7	41	69	58951	46	33	6	99	73

	Fwd Band- I	Forwarded	Avg Read	Avg Write	Read Band- I	Write Band- I	Total Band- I	Read I	Write
Idx	Width(KiB/s)	IOPS	Latency(ms)	Latency(ms)	Width(KiB/s)	Width(KiB/s)	Width(KiB/s)	IOPS	IOPS
0	0.00	0.00	9.09	0.54	1580.95	2663.77	4244.73	1.9	7.
1	0.00	0.00	0.15	0.00	0.00	0.00	0.00	0.0	0.
2	0.00	0.00	8,81	0.66	1624,55	2794.02	4418.57	2.1	7.
3	0.00	0.00	0.11	0.00	0.00	0.00	0.00	0.0	0.
4	0.00	0.00	8.59	0.55	1761.36	2761.23	4522.59	2.2	8.
5	0.00	0.00	0.09	0.00	0.00	0.00	0.00	0.0	0.
6	0.00	0.00	8.69	0.53	1595.09	2767.30	4362.38	2.1	7.
7	0.00	0.00	0.08	0.00	0.00	0.00	0.00	0.0	0.
8	0.00	0.00	8.89	0,54	1988.54	3230.76	5219.30	2.4	8
9	0.00	0.00	0.10	0.00	0.00	0.00	0.00	0.0	0
10	0.00	0.00	8.47	0.54	1929.38	2898.30	4827.68	2.3	7.
11	0 00	0 00	0 00	0.00	0.00	0 00	0 00	0 0	0

- At ORNL we use Nagios
 - Provides a dashboard for system health
- Monitoring
 - OK, warning, critical
- Alerting
 - Business hours
 - non-business hours
 - never page





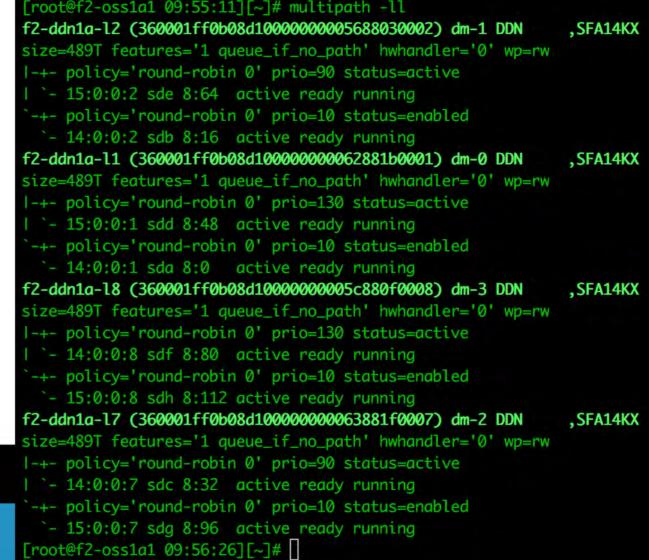
- block-device tuning checks
 - Need to make sure that the IO scheduler, nr_requests, timeouts, etc are tuned correctly.

 [root@f2-oss1al 09:21:02][~]# ipmitool sdr
 - These can be lost after an upgrade.
- Mounted devices check:
 - Make sure that all of your OSTs are mounted
 - sounds ridiculous, but this can happen at 2AM
- Server health (memory/processor/fan/powersupply)
 - Many hardware vendors provides tools already
 - OpenManage (Dell), iLO (HP), etc...
 - ipmitool sdr
 - Hundreds of sensors available



_ ·		,		/
[root@f2-oss1a1	09	:21:02][~]#	ipmitool	sdr
		29 degrees		
Temp		29 degrees	C	ok
Inlet Temp		15 degrees	C	ok
DIMM PG		0x00	1	ok
NDC PG	1	0x00	1	ok
PS1 PG FAIL	1	0x00	- 1	ok
PS2 PG FAIL	1	0x00		ok
BPØ PG		0×00		ok
BP1 PG		0×00	1	ok
1.8V SW PG			1	ok
2.5V SW PG	1	0×00	1	ok
5V SW PG	1	0×00	1	ok
PVNN SW PG	1	0×00	1	ok
VSB11 SW PG	1	0×00	1	ok
VSBM SW PG	1	0×00	1	ok
3.3V B PG	1	0x00	1	ok
MEM012 VDDQ PG		0x00	1	ok
MEM012 VPP PG		0×00	1	ok
MEM012 VTT PG		0×00	1	ok
MEM345 VDDQ PG		0×00		ok
MEM345 VPP PG		0×00		ok
MEM345 VTT PG		0×00		ok
VCCIO PG		0×00	1	ok
VCORE PG		0×00	1	ok
FIVR PG		0x00	1	ok
MEM012 VDDQ PG		0×00	1	ok
MEM012 VPP PG		0×00	1	ok
MEM012 VTT PG		0×00	1	ok
MEM345 VDDQ PG	1	0x00	1	ok
MEM345 VPP PG	1	0×00	1	ok
MEM345 VTT PG		0×00	1	ok
VCCIO PG		0×00	1	ok
VCORE PG		0×00		ok
FIVR PG		0×00		ok
Fan1A		5280 RPM		ok
D 40		1000 0000		

- Critical services monitoring:
 - Want to make sure that services that are required for system operation are "running"
 - Examples: srp_daemon/opensmd/crond/postfix
 - Simple script that parses `systemctl status <service>` output
- Multipath health
 - Script that parses `multipath –ll`
 - 2 paths: Healthy,
 - 1 path: warning,
 - O paths: critical





- Host IB health
 - Network link health (lane count and speed)
 - Check for card->PCI bus link health
 - Check counters changes over time
 - Symbol errors
 - LinkDownedCounter
 - VL15 Dropped

[root@f2-oss1a1 10:37:56][lustre]# lspci -s 3b:00.0 -vvv

3b:00.0 Infiniband controller: Mellanox Technologies MT27700 Family [ConnectX-4] Subsystem: Mellanox Technologies Device 0014 Control: I/O- Mem+ BusMaster+ SpecCycle- MemWINV- VGASnoop- ParErr- Stepping- SERR- FastB2B- DisINTx+ Status: Cap+ 66MHz- UDF- FastB2B- ParErr- DEVSEL=fast >TAbort- <TAbort- <TAbort- <SERR- <PERR- INTx-Latency: 0, Cache Line Size: 32 bytes Interrupt: pin A routed to IRQ 53 NUMA node: 0 Region 0: Memory at ae000000 (64-bit, prefetchable) [size=32M] Expansion ROM at ab000000 [disabled] [size=1M] Capabilities: [60] Express (v2) Endpoint, MSI 00 DevCap: MaxPayload 512 bytes, PhantFunc 0, Latency LOs unlimited, L1 unlimited ExtTag+ AttnBtn- AttnInd- PwrInd- RBE+ FLReset+ SlotPowerLimit 75.000W DevCtl: Report errors: Correctable- Non-Fatal+ Fatal+ Unsupported+ RlxdOrd+ ExtTag+ PhantFunc- AuxPwr- NoSnoop+ FLReset-MaxPayload 256 bytes, MaxReadReg 512 bytes DevSta: CorrErr+ UncorrErr- FatalErr- UnsuppReq+ AuxPwr- TransPend-LnkCap: Port #0, Speed 8GT/s, Width x16, ASPM not supported, Exit Latency L0s unlimited, L1 unlimited ClockPM- Surprise- LLActRep- BwNot- ASPMOptComp+ LnkCtl: ASPM Disabled; RCB 64 bytes Disabled- CommClk-ExtSynch- ClockPM- AutWidDis- BWInt- AutBWInt-LnkSta: Speed 8GT/s, Width x16, TrErr- Train- SlotClk+ DLActive- BWMgmt- ABWMgmt-



<pre>[root@f2-oss1a1 10:40:01][lustre]# i CA 'mlx5_0'</pre>	bstat
CA type: MT4115	
Number of ports: 1	
Firmware version: 12.18.1000)
Hardware version: 0	
Node GUID: 0x506b4b03003956b)e
System image GUID: 0x506b4b0	
Port 1:	
State: Active	
Physical state: Link	dln
Rate: 100	(OP
Base lid: 1	
LMC: 0	
SM lid: 1	
Capability mask: 0x2	6510840
Port GUID: 0x506b4b0	
Link layer: InfiniBa	
[root@f2-oss1a1 10:41:16][lustre]# p	
<pre># Port counters: Lid 1 port 1 (CapMa</pre>	sk: 0x5A00)
PortSelect:1	
CounterSelect:0x0	000
CounterSelect:0x0 SymbolErrorCounter:0	000
CounterSelect:0x0 SymbolErrorCounter:0 LinkErrorRecoveryCounter:0	000
CounterSelect:0x0 SymbolErrorCounter:0 LinkErrorRecoveryCounter:0 LinkDownedCounter:0	000
CounterSelect:0x0 SymbolErrorCounter:0 LinkErrorRecoveryCounter:0	000
CounterSelect:0x0 SymbolErrorCounter:0 LinkErrorRecoveryCounter:0 LinkDownedCounter:0	000
CounterSelect:0x0 SymbolErrorCounter:0 LinkErrorRecoveryCounter:0 LinkDownedCounter:0 PortRcvErrors:0 PortRcvRemotePhysicalErrors:0 PortRcvSwitchRelayErrors:0	000
CounterSelect:0x0 SymbolErrorCounter:0 LinkErrorRecoveryCounter:0 LinkDownedCounter:0 PortRcvErrors:0 PortRcvRemotePhysicalErrors:0	000
CounterSelect:0x0 SymbolErrorCounter:0 LinkErrorRecoveryCounter:0 LinkDownedCounter:0 PortRcvErrors:0 PortRcvRemotePhysicalErrors:0 PortRcvSwitchRelayErrors:0	000
CounterSelect:0x0 SymbolErrorCounter:0 LinkErrorRecoveryCounter:0 LinkDownedCounter:0 PortRcvErrors:0 PortRcvRemotePhysicalErrors:0 PortRcvSwitchRelayErrors:0 PortXmitDiscards:0	000
CounterSelect:0x0 SymbolErrorCounter:0 LinkErrorRecoveryCounter:0 LinkDownedCounter:0 PortRcvErrors:0 PortRcvRemotePhysicalErrors:0 PortRcvSwitchRelayErrors:0 PortXmitDiscards:0 PortXmitDiscards:0	
CounterSelect:.0x0SymbolErrorCounter:.0LinkErrorRecoveryCounter:.0LinkDownedCounter:.0PortRcvErrors:.0PortRcvRemotePhysicalErrors:.0PortRcvSwitchRelayErrors:.0PortXmitDiscards:.0PortXmitConstraintErrors:.0PortRcvConstraintErrors:.0	
CounterSelect:.0x0SymbolErrorCounter:.0LinkErrorRecoveryCounter:.0LinkDownedCounter:.0PortRcvErrors:.0PortRcvRemotePhysicalErrors:.0PortRcvSwitchRelayErrors:.0PortXmitDiscards:.0PortXmitConstraintErrors:.0PortRcvConstraintErrors:.0CounterSelect2:.0x0	
CounterSelect:	
CounterSelect:.0x0SymbolErrorCounter:.0LinkErrorRecoveryCounter:.0LinkDownedCounter:.0PortRcvErrors:.0PortRcvRemotePhysicalErrors:.0PortRcvSwitchRelayErrors:.0PortXmitDiscards:.0PortXmitConstraintErrors:.0PortRcvConstraintErrors:.0CounterSelect2:.0x0LocalLinkIntegrityErrors:.0ExcessiveBufferOverrunErrors:.0	0
CounterSelect:.0x0SymbolErrorCounter:.0LinkErrorRecoveryCounter:.0LinkDownedCounter:.0PortRcvErrors:.0PortRcvRemotePhysicalErrors:.0PortRcvSwitchRelayErrors:.0PortXmitDiscards:.0PortXmitConstraintErrors:.0PortRcvConstraintErrors:.0CounterSelect2:.0x0LocalLinkIntegrityErrors:.0VL15Dropped:.0	0 4967295
CounterSelect:.0x0SymbolErrorCounter:.0LinkErrorRecoveryCounter:.0LinkDownedCounter:.0PortRcvErrors:.0PortRcvRemotePhysicalErrors:.0PortRcvSwitchRelayErrors:.0PortXmitDiscards:.0PortXmitConstraintErrors:.0PortRcvConstraintErrors:.0PortRcvConstraintErrors:.0CounterSelect2:.0x0LocalLinkIntegrityErrors:.0VL15Dropped:.0PortXmitData:.429	0 4967295 4967295
CounterSelect:.0x0SymbolErrorCounter:.0LinkErrorRecoveryCounter:.0LinkDownedCounter:.0PortRcvErrors:.0PortRcvRemotePhysicalErrors:.0PortRcvSwitchRelayErrors:.0PortXmitDiscards:.0PortXmitConstraintErrors:.0PortRcvConstraintErrors:.0CounterSelect2:.0x0LocalLinkIntegrityErrors:.0VL15Dropped:.0PortXmitData:.429PortRcvData:.429	0 4967295 4967295 4967295
CounterSelect:.0x00SymbolErrorCounter:.0LinkErrorRecoveryCounter:.0LinkDownedCounter:.0PortRcvErrors:.0PortRcvRemotePhysicalErrors:.0PortRcvSwitchRelayErrors:.0PortXmitDiscards:.0PortXmitConstraintErrors:.0PortRcvConstraintErrors:.0CounterSelect2:.0x00LocalLinkIntegrityErrors:.0VL15Dropped:.0PortXmitData:.429PortRcvData:.429PortXmitPkts:.429	0 4967295 4967295 4967295 4967295

- Switch-to-switch IB health
 - `ibdiagnet` is insufficient for finding all switch-to-switch IB link issues
 - It will help you find unhealthy links
 - Links can go dark and will not be detected
 - Down IB links can cause performance issues
 - If using AR, they can even make the network re-route which causes unavailability for a small time
 - Non-symmetric routes can cause ~3-5% performance drop
 - Script that knows IB network topology and checks for it to be sane
 - Knows that the switch cables are connected to the correct port on the correct switch
 - Impacts network routing
 - Knows that the host cables are connected to the correct port on the correct switch
 - If improperly cabled can impact FGR
 - Check link speed
 - Check link width



- Lnet_stats
 - Monitor changes in `Inetctl stats show` to show LNET congestion or errors
 - Set threshold to report on changes in backlogged messages "msg_alloc" – Example 30000
 - Can set a threshold for downed routes, dropped messages, etc...
- Lustre_health
 - Simple script that checks the status of `lctl get_param -n health_check`
 - Tells you if a OST is mounted read-only, is slow, corrupt, etc...
- Ls timer
 - Inside of each cluster network (if routed), check to make sure that you can `ls` inside of a lustre directory within a certain timeout
 - Will tell you if lustre is being slow or not
 - Helps to get in front of users complaining



15-minute break



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When Things Go Wrong

- Lustre Recovery
- Gathering debug information
- Network debugging
- Repairing file system issues



Lustre Recovery

- Lustre's recovery mechanism is designed to deal with node/network failures and keep the file system running in a consistent state
- Some of the failures it is designed to handle are:
 - 1. Client failure
 - 2. MDT failure
 - 3. OST failure
- MDS and OSS failures require methods to recover or replay outstanding I/O requests from clients



Client Failure - Detection

- It is important to detect client failure early so that remaining clients can continue accessing the file system
- Two main ways to detect client failure:
 - 1. Client fails to respond to a blocking lock callback from the Distributed Lock Manager (DLM)
 - 2. Client fails to "ping" server in a long period of time
- These conditions may occur even if the hardware itself has not actually failed (e.g. – network link failure), but it is still treated the same



Client Failure - Recovery

- When a client failure is detected, Lustre tries to ensure that other clients can continue working
 - Can't afford to have one or more clients waiting to perform I/O while they are trying to acquire a lock held by a dead client
- When a client is evicted:
 - All client locks are invalidated
 - All cached inodes on client are invalidated
 - All cached data on client is flushed
- When client recovers, it may reconnect to the file system and continue operations



MDT Failure - Detection

- Clients may detect MDT failure by timeouts of in-flight requests or from Imperative Recovery
 - Client MDC will attempt to connect to failover node if configured
 - Only clients connected during the failure are permitted to reconnect during the recovery window
- Client state will need to be communicated to MDT once connection is reestablished



MDT Failure - Recovery

- Lustre uses the Metadata Replay protocol to ensure that MDS can re-acquire necessary state information from client transactions that have not been committed to disk
- The protocol uses transaction numbers to ensure operations are replayed in the correct order
- Clients also communicate existing lock state to MDS



OST Failure - Detection

- If an OST fails to respond to a client in a timely manner, the corresponding OSC on the client will treat the OST as having failed
 - Outstanding I/O requests will block until the OST has recovered
 - OSC will try to reconnect to OST through a failover OSS node (if one has been configured)
- Same logic applies if the "client" is the MDS
 - MDS will note that OST is unavailable and skip it when assigning objects to new files



OST Failure - Recovery

- OSC-to-OST recovery protocol is the same as the MDC-to-MDT Metadata Replay protocol
 - Bulk writes usually have been committed to disk so server just needs to reconstruct the reply
 - For other cases, normal replay/resend handling is done
 - Client still has copy of data until it receives acknowledgement
- When OST is in recovery mode, all new client connections are refused until the recovery finishes
 - Recovery finishes when all previously-connected clients have replayed transactions, or a client times out



Metadata Replay Protocol

- Every client request contains a unique, monotonically increasing XID to track order of requests
- Each request processed by server is assigned a unique, increasing Transaction Number (TN)
 - Reply to client's request contains TN for the request along with the last committed TN
- Server maintains last_rcvd file with list of connected clients
- During recovery
 - Request with only XID \rightarrow resend
 - Request with TN \rightarrow replay



Viewing Recovery Status

- To view recovery status of all OSTs
 lctl get_param obdfilter.*.recovery_status
- To view recovery status of MDTs

lctl get_param mdt.*.recovery_status

- Example output:
 - status: COMPLETE
 recovery_start: 1553204504
 recovery_duration: 0
 completed_clients: 1/1
 replayed_requests: 0
 last_transno: 94574301709



Aborting Recovery

- In some cases, it may be known that recovery will not complete properly, or perhaps recovery is not really necessary
 - Previously connected client may currently be down
 - File system was brought down cleanly, but there was an idle client connected at the time
- Recovery can be aborted in two ways:

```
mount -t lustre -o abort_recov <dev> <mnt_point>
```

```
lctl --device <dev_num> abort_recovery
```



Gathering Debug Information

- When something goes wrong with Lustre, there are several ways to grab useful information
- Some of these methods are useful for sys admins, and others are primarily of use to developers
- Sources of debug information include:
 - Syslog / dmesg
 - Lustre internal debug logs
 - Crash dumps
 - Debugfs
 - Wireshark



Syslog / dmesg

- Things to look for in log messages:
 - Lustre / LustreError / Lnet / LBUG
 - rc -30 (EROFS)
 - Timeouts / evictions
 - Messages that contain NIDs (<u>10.1.2.3@o2ib</u>, etc.)
- It's impossible to enumerate all the Lustre errors you might see, so let Google be your friend
- Sometimes general pattern of messages can be just as useful (or perhaps more useful) than the content of the messages



Lustre Internal Debug Log

- Lustre maintains an internal circular debug buffer
- A debug mask is used to control what info gets logged
 - Query using "lctl get_param debug"
 - Setusing "lctl set_param debug=<mask>"
 - Can also be set using "sysctl lnet.debug"
- Size of the buffer can be modified using lctl set_param debug_mb=<size>
- Contents of buffer can be dumped to a file using lctl debug_kernel <filename>
- See Lustre manual for info about debug mask options



Crash Dumps

- Use kdump (via kexec) to capture kernel info when LBUG is encountered
 - 1. Set kernel to panic on LBUG

lctl set_param panic_on_lbug=1

- 2. Install kexec-tools package
- 3. Add the following parameter to the kernel boot options:

crashkernel=<size> (or "auto")

- 4. Modify / etc/kdump.conf if desired
 - For example, send crash dumps over the network to another host
- 5. Start the kdump service

systemctl start kdump

• Use a program like crash to analyze output



debugfs

- When using Idiskfs for backend Lustre storage, you can inspect the contents of the file system in two ways:
 - 1. Mount the device with "-t ldiskfs" instead of "-t lustre"
 - 2. Use the debugfs command
- One benefit of using debugfs is that you can view the contents while Lustre is up and running
 - In that case, it is best to use "debugfs -c" so that the device is opened in read-only mode
- Even if there are no problems, spending some time looking at the file layout can provide some insight into how Lustre works



Network Debugging

- Many Lustre issues can ultimately be traced back to network connectivity problems
 - Disruption of client-server communication leads to timeouts or dropped requests
 - Clients see this as a server failure
 - Servers see it as a client failure and evict clients
- Error messages might not make the issue obvious
 - Client syslog message may complain about being unable to process config from MDS, but the real reason is that it can't even contact the MDS
 - Sporadic network problems make debugging even harder



Initial Troubleshooting

- Is the firewall enabled?
- Does every node have the proper NID configured on the correct interface?
- If LNet routing is used, does the client and server have the correct routes?
- Do any nodes have duplicate IP addresses?
- Can you ping between nodes? Both ways?
- Can you lctl ping between nodes? Both ways?
- Do the servers have MDTs/OSTs mounted?



Infiniband Issues

- Debugging Infiniband issues can get complex, but there are some simple steps that often lead to results
 - Is IPoIB configured?
 - Is the installed version of Lustre built against the correct IB stack (inkernel vs. MOFED, version, etc.)?
 - Is the IB firmware too old? Too new?
 - Do IB bandwidth tests give expected results?
 - Do IB HBA counters show any errors? What about counters on the IB switch?
 - Does output from ibnetdiscover, ibstat, etc. match what you expect?



Lnet selftest

- Lnet selftest is a useful tool for testing connectivity and measuring network performance
- Can be used to test pairs of nodes or entire clusters
- To run Lnet selftest:
 - Load lnet_selftest kernel module on all nodes
 - Use lst command to add groups of clients and servers
 - Use lst command to specify type of test to run
 - Initiate test from any host on the fabric



Example: Lnet selftest

export LST_SESSION=\$\$ echo LST_SESSION=\$LST_SESSION

Ist new_session io_test Ist add_group clients 10.10.20.31@o2ib0 Ist add_group servers 10.10.1.7@o2ib0 10.10.1.8@o2ib0 Ist add_batch bulk Ist add_test --batch bulk --concurrency=8 --distribute 1:2 --from clients \ --to servers brw write size=1M Ist run bulk Ist stat servers & sleep 30; kill \$! Ist end_session



Repairing File System Issues

- Sometimes file system data structures can get into an inconsistent state
- Causes can include:
 - Power failures
 - Hardware failures
 - Software bugs
- Inconsistency could be with Lustre's internal data or with the data structures of the backend ldiskfs/zfs file systems used on the MDTs/OSTs
- Each layer has its own tools to deal with the problem

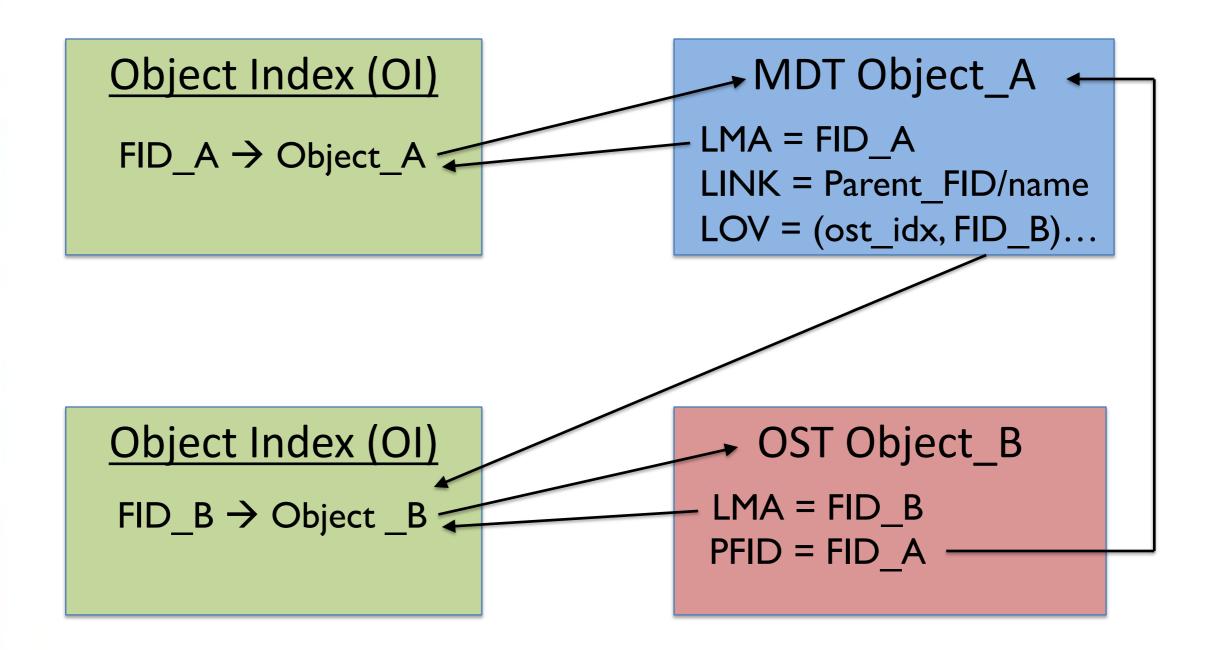


LFSCK (Lustre File System Checker)

- Lustre provides a tool for checking the consistency of its internal state and repairing any problems
- Prior to Lustre 2.3, performing a full file system check was sloooooow and painful
 - Had to take Lustre offline to generate the needed databases of inode information
 - Best bet was just to run e2fsck on underlying ldiskfs file system and hope it fixed enough of the problems
- LFSCK has been re-engineered to run with the file system online (and in use)



File Identifiers and Objects





LFSCK (Phase 1)

- Maintain consistency of Object Index on MDT
- Iterate through all objects on the OSD
 - Make sure inode number in OI matches with FID from inode's LMA xattr
- Can be triggered manually or automatically
- Maintains checkpoint file (scrub_status) on MDT
 - Allows restart if scan is interrupted
 - Contains stats about current scan
- Supports rate limiting



LFSCK (Phase 1.5)

- Maintain consistency between the FID-in-Dirent info and LMA/LINK xattrs in objects
- Iterate through each object on OSD
- If it is a directory, check each file entry
 - Compare FID listed in dirent with LMA xattr of inode
 - Compare file name from dirent with name from inode's LINK xattr
 - Compare FID from LINK xattr with FID of parent directory
- Supports checkpoint restart (Ifsck_namespace) and rate limiting
- This check is not automatically triggered



LFSCK (Phase 2)

- MDT-OST consistency checking
- MDT object for a file contains list of child OST objects
- Child OST object contains FID for parent MDT object
- Check 4 different cases:
 - Dangling reference mdt_obj1 points to ost_obj1, but ost_obj1 doesn't exist or doesn't have PFID xattr
 - Mismatched reference mdt_obj1 points to ost_obj1, but ost_obj1 points to mdt_obj2. mdt_obj2 doesn't exist or recognize ost_obj1 as child.
 - Multiple references mdt_obj1 and mdt_obj2 both point to ost_obj1
 - Unreferenced object ost_obj1 points to mdt_obj1, but mdt_obj1 doesn't exist or recognize ost_obj1 as child. No other mdt_obj points to ost_obj1.



LFSCK (Phase 2)

- Fixes ownership inconsistency between MDT and OST objects (MDT ownership takes precedence)
- Will track errors, and if threshold is reached, will trigger full lfsck for file system
- Supports checkpoint restart and rate limiting



LFSCK (Phase 3)

- Implements MDT-MDT consistency check for DNE
- Similar to MDT-OST consistency check in many ways, but also more complicated
- Too many cases to list here
 - Check <u>http://wiki.lustre.org</u> for design docs
- Supports checkpoint restart and rate limiting



Running LFSCK

- Full file system check is initiated via
 lctl lfsck_start-M \${MDT0} -A -t all -r
- The -t option is used to specify which checks to run
 - scrub Run OI scrub
 - namespace FID-in-Dirent, LinkEA consistency
 - Iayout MDT-OST object consistency
- Other useful options
 - -n | --dryrun
 - -c | --create_ostobj
 - -C | --create_mdtobj
 - -o | --orphan



Example: Running LFSCK

[root@haven-mds1 ~]# lctl lfsck_start -M haven-MDT0000 -A -t all -r

```
[root@haven-mds1 ~]# lctl lfsck_query -M haven-MDT0000
layout_mdts_init: 0
layout_mdts_scanning-phase1: 1
layout_mdts_scanning-phase2: 0
```

layout_osts_scanning-phase1: 30
layout_osts_scanning-phase2: 12

namespace_mdts_init: 0
namespace_mdts_scanning-phase1: 1
namespace_mdts_scanning-phase2: 0



Repairing Idiskfs corruption

- Since Idiskfs is based on ext4, journaling helps keep the file system in a consistent state
- If a problem occurs that cannot be fixed by the journal, it will be necessary to run e2fsck
 - One possible symptom of this is when the logs contain "-30" (EROFS) errors
 - Only need to run e2fsck on the device(s) that contain errors
- General procedure:
 - 1. Replay journal
 - 2. Run e2fsck in non-fixing mode
 - 3. Run e2fsck to fix problems



Example: Running e2fsck

NOTE: Always use latest e2fsprogs from Whamcloud https://downloads.whamcloud.com/public/e2fsprogs/latest/

Unmount affected device
root# umount /mnt/ost

If possible, use logger to capture output
root# script /tmp/e2fsck.sda

Replay journal
root# mount -t ldiskfs /dev/sda /mnt/ost
root# umount /mnt/ost



Example: Running e2fsck (cont.)

```
# Run e2fsck in non-fixing mode
root# e2fsck -fn /dev/sda
...[output]...
```

```
# Fix the errors
root# e2fsck -fp /dev/sda
...[output]...
```

Might need to follow-up with LFSCK if there are lots of problems



ZFS maintenance

- ZFS handles consistency issues differently from ldiskfs
- Admins should periodically scrub zpools
 - Can be done while zpool is online and Lustre is running
 - Causes I/O to disk which could have some affect on the file system
 - Recommended interval = 1 month (?)
- Example:

```
zpool scrub <pool_name>
```

 Can reduce impact from scrub by adjusting sysctl parameter vfs.zfs.scrub_delay



Other Useful Admin Info

- Striping Considerations
- OST allocation (Round-robin vs. Weighted)
- Advanced file layouts
 - Progressive File Layout (PFL)
 - Data on MDT (DoM)



Striping Considerations

- Basic file striping is pretty straightforward
 - Most of the time, just choose a stripe count
 - Sometimes you might adjust the stripe size
 - Other options probably used even less
- For user, striping is usually about performance
 Knowledge of application IO pattern

Customized striping parameters

$\mathbf{\Psi}$

Less contention, better IO performance

• But admins have additional concerns...



Default stripe count

- Choosing the default stripe count for a file system can be a tricky proposition
 - Too low \rightarrow Fill up OSTs with large files
 - Too high \rightarrow Consume more inodes on OSTs than needed
 - Progressive File Layouts can help with this
- Choice of default stripe count might also affect how you choose to format the MDTs/OSTs
- In any case, it's a good idea to have some general guidelines for users
 - Ex At least 1 stripe for every 100 GB of file space



Improperly striped files

- Whatever striping guidelines you choose, users still won't listen...
- May need to track down large files with small stripe counts that are filling up OSTs
- Options:
 - 1. Ifs find (could take a while)
 - 2. Robinhood (if you already have this tool)
 - 3. OST usage distribution (quick, but limited)
- The last option is handy, but sometimes requires a little work



Searching for Improperly Striped Files

- Look at distribution of OST usage
 - Run "Ifs df <filesystem>| sort -nk 5"
 - Look for anomalies at the tail end
- Find the user(s) with the most usage on OST
 - Run "lfs quota -l <ost_idx> -u <user> <filesystem>" command for each user
 - Look for one or more users with abnormally high usage
 - These are your initial candidates for investigation
- Try to locate the offending files



Example

haven-OST000e_UUID x x x 40% /lustre/haven[OST:14] haven-OST001a_UUID x x x 41% /lustre/haven[OST:26] haven-OST0017_UUID x x x 41% /lustre/haven[OST:23] ...<snip>...

haven-OST0019_UUID x x x 51% /lustre/haven[OST:25] haven-OST0005_UUID x x x 52% /lustre/haven[OST:5] haven-OST0008_UUID x x x 52% /lustre/haven[OST:8] haven-OST0013_UUID x x x 53% /lustre/haven[OST:19] haven-OST001b_UUID x x x 65% /lustre/haven[OST:27] filesystem summary: x x x 46% /lustre/haven



Inode Calculations

- Default stripe count and average file size are important factors for planning a new file system
 - These factors help determine the number of inodes needed on MDTs/OSTs which in turn can affect formatting options and device size requirements
- Number of MDT inodes:
 - num_osts * ost_size / avg_file_size
 - Recommend doubling this to allow for future expansion or smaller than expected file size
- Number of OST inodes:
 - num_mds_inodes * default_stripe_count / num_osts
 - Recommend 2x-4x padding



Inode Calculations (cont.)

- ZFS has variable number of inodes
 - MDT still needs enough space to allow about 4KB per inode
- Idiskfs creates fixed number of inodes during format
- Defaults for Lustre 2.10:
 - inode size = 1KB
 - MDT will have 1 inode for every 2.5KB
 - OST will have 1 inode for every 1MB (if OST size > 8TB)
- Can alter inode ratios by adding option to mkfs.lustre:
 - --makefsoptions="-i <bytes-per-inode>"
- Adjust this option to get desired number of inodes



Inode Disparity

- Primary goal of these calculations is to have parity among MDT and OST inode counts
 - Ideally, inode and space usage track each other
- Disparity can show up in non-obvious ways # Ifs df -i /Ifs01

UUID	Inodes	IUsed	lfree l	IUse% Mounted on
MDT0000	240228761	6 46560885	235572673	1 2% /share/lfs01[MDT:0]
OST0001	24117248	22883788	1233460	95% /share/lfs01[OST:1]
OST0003	24117248	22903308	1213940	95% /share/lfs01[OST:3]
OST0004	24117248	22895442	1221806	95% /share/lfs01[OST:4]
OST0006	24117248	22890201	1227047	95% /share/lfs01[OST:6]

summary: 51457138 46560885 4896253 90% /share/lfs01



OST Object Allocation

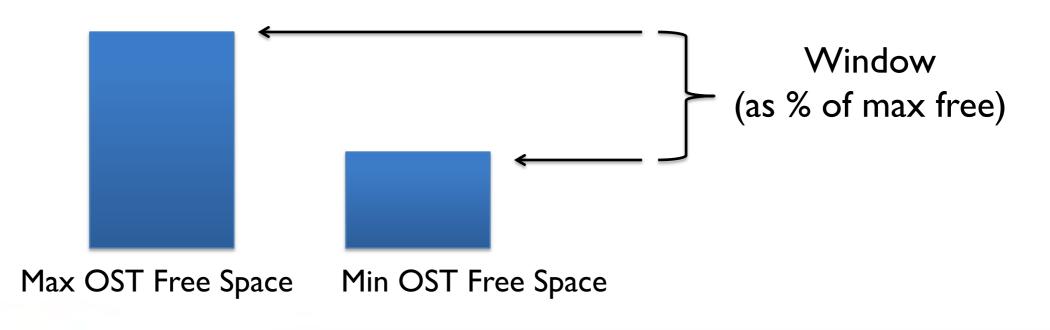
 When a new file is created, Lustre allocates objects on OSTs according to desired stripe count

-bash-4.2\$ lfs get testfile	tstripe testfile	2			
Imm_stripe_cou	nt: 4				
lmm_stripe_size: 1048576					
lmm_pattern: 1					
Imm_layout_gen: 0					
Imm_stripe_offset: 8					
obdid	objid	objid			group
8	231338244	Oxdc	9f104		0
39	20273590	0x13559b6	0		
30	20441490	0x137e992	0		
38	20549867	0x13990eb	0		



OST Allocators

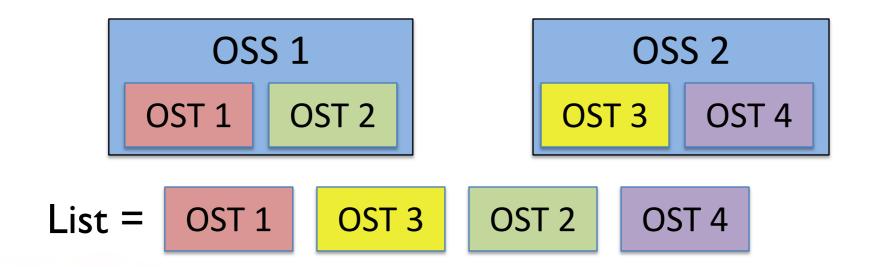
- How does Lustre decide which OSTs to assign to a file?
- Two different allocators
 - Round-robin
 - Weighted
- Choice based on how "balanced" usage is (as defined by the target window)





Round-Robin Allocator

- Round-robin allocator is used if the OST usage is balanced (i.e. – all OST free space falls within target window)
- OSTs are assigned sequentially from an internal list
- List is not necessarily sequential with regards to OST index
 - Accounts for things like OSTs being on different nodes





Weighted Allocator

- Weighted allocator is used when OST usage is not balanced
- OSTs are assigned a weight based on the amount of free space and their location
 - Emptier OSTs have a higher weight and are more likely to be selected
- Algorithm makes random selection based on weights
 - Even OSTs with the least free space still have some chance of being selected
- The goal is to divert more I/O to OSTs with the most free space while still utilizing other OSTs to some extent



Adjusting Allocator

- Admins have some control over which allocator is used and how the Weighted allocator assigns weights
- Control size of window used to determine if OST usage is balanced
 - /proc/fs/lustre/lov/<name>-MDT0000-mdtlov/qos_threshold_rr
 - Default value is 17%
 - If set to 100%, round-robin is always used
- Control how much weight is affected by free space
 - /proc/fs/lustre/lov/<name>-MDT0000-mdtlov/qos_prio_free
 - Default value is 91%
 - If set to 100%, weights are based solely on free space



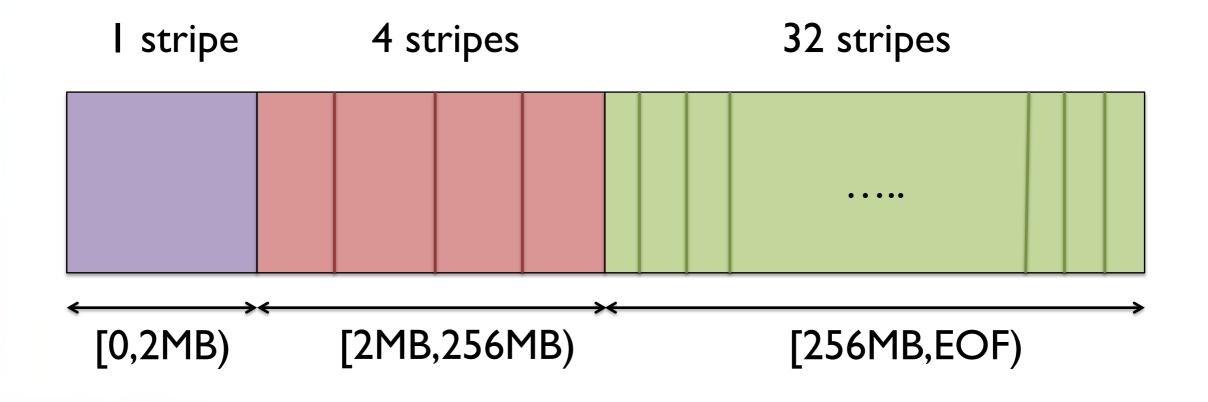
Advanced File Layouts

- Recent versions of Lustre have added some features that provide more options beyond current basic layout
- 1. Progressive File Layout
 - Provides ability to adjust file layout as the size of the file grows.
 - Essentially creates different basic layouts for different sections of a file
- 2. Data on MDT
 - Store some (or possibly all) file contents on the MDT itself



Progressive File Layout (PFL)

- Introduced in Lustre 2.10
- A PFL file is essentially an array of basic layouts (components) that cover different non-overlapping sections of a file





PFL Benefits

- Fine-grain control of layout could provide performance improvements
- File layout can be adapted on-the-fly
 - Only need to define initial component
 - Add components when needed
 - Don't use more OST inodes than necessary
- Choose a default PFL for all users that gradually increases stripe count as the file size increases
 - No more full OSTs! (maybe...)
- Underlying composite layout structure forms basis for other layout options



PFL Examples

Create PFL for previous figure Ifs setstripe -E 2M -c 1 -E 256M -c 4 -E -1 -c 32 <file>

Create starting layout, then add component Ifs setstripe -E 2M -c 1 -E 256M -c 4 <file> Ifs setstripe --component-add -E -1 -c 32 <file>

Display all components of file
lfs getstripe <file>

NOTE: Will only see OST objects for instantiated components



Data on MDT (DoM)

- Introduced in Lustre 2.11
- Designed to improve file I/O by placing small files (or the first part of a larger file) directly on MDT
 - Helps eliminate extra RPCs to OSTs
 - Advantageous if MDT storage is faster than OST storage
- This is a special case of PFL in which the first component has a single stripe that resides on the MDT
- Example:

lfs setstripe -E 1M -L mdt -E 256M -c 4 -E EOF -c 10 <file>



DoM Settings

- Some care must be taken when allowing users to place data directly on the MDT
- Admins can limit the size of the file's first stripe that resides on the MDT
- Controlled via dom_stripesize parameter (default=1MB, disabled=0):

```
# Query value
```

lctl get_param lod.*MDT0000*.dom_stripesize

Set value temporarily

lctl set_param lod.*MDT0000*.dom_stripesize=<value>

Set value permanently

lctl conf_param <fsname>-MDT0000.lod.dom_stripesize=<value>



Panel Session



Questions?



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