

A Vision of Storage for Exascale Computing

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Fast Forward Storage & IO Project Goals

- Make Exascale storage a tool of the Scientist
 - Tractable data management
 - Comprehensive interaction
 - Move compute to data or data to compute as appropriate
- Overcome today's IO limits
 - Multi-petabyte datasets
 - Explosive growth of metadata
 - Horizontal scaling & jitter
- Support unprecedented fault tolerance
 - Deterministic interactions with failing hardware & software
 - Fast & scalable recovery
 - Enable multiple redundancy & integrity schemes



Fast Forward I/O Architecture



I/O forwarding from compute to IO nodes



Workflow: Simulation + In-transit Analysis





Workflow: Pre-stage to Burst Buffer



Scientist Workstation

 Pre-stage triggered when BB resources released by previous workflow session



Workflow: Start Session



 Previous session may still be persisting data from BB to global storage



Workflow: Dump Timestep





Workflow: In-transit Analysis



Raw timestep may be discarded



Workflow: Persist to Global Storage



Scientist Workstation

Analysis output saved to global storage



Workflow: Browse



 Insufficient bandwidth for brute-force query or index build



Workflow: Analysis Shipping



Scientist Workstation

- Ship index build / query to storage cluster
- Full streaming bandwidth available
- Query results returned to workstation



Stackable components

- Application I/O APIs
 - Multiple domain-specific API styles & schemas
 - HDF5+extensions & Graph Computation libraries
- I/O forwarding
 - Keeps top level APIs on Compute Nodes when IOD runs on the Burst Buffer
- I/O Dispatcher (IOD)
 - Impedance match application I/O to storage capabilities
 - Semantic resharding
 - Burst Buffer management
- DAOS-HA
 - High-availability scalable object storage
 - Follow-on project from Fast Forward
- DAOS Containers
 - Virtualized shared-nothing object storage
 - Unpolluted storage system namespace





*other names and brands may be claimed by others



I/O Stack Configurations



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Ubiquitous NVRAM

- O(1TB) compute node-local storage
- Instant-on
 - 0 power standby
- Load-store byte-granular access
 - Invites Distributed Persistent Memory programming models
 - Order of magnitude larger in-core working sets
- Storage fully leverages fabric

	Disk	Edge BB	NVRAM
Checkpoint / Search	1 hour	6 minutes	6 seconds
Capacity (# datasets)	30	3-5	10-30



Scheduling Persistent Memory



- Workflow Session 4 ready to run
- Data not local
- Migrate
- Workflow Session 4 started

- Issues
 - Space at destination
 - Comms Interference



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Persistent Memory v. Storage

- Persistent Memory is fast but it's...
 - Local to the process using it
 - Inaccessible on node failure
 - Fixed schema
- Storage may be slower but it's...
 - Globally accessible
 - Consistent & durable
 - Snapshotable / Cloneable / Migrateable
- APIs required to...
 - − Convert PM ⇔ Storage
 - Persist / Instantiate Distributed Persistent Memory images
 - PM schema conversion
 - Support workflow scheduler integration
 - Data-aware process instantiation
 - Process-aware data migration





DAOS-M

- Client & Server OS bypass
- Connectionless
 - Peer-to-peer connectivity = ~100x client/server
 - Heavyweight security / ownership checks once on container open
- Memory VOSD
 - PM programming model
 - No block I/O stack latency
 - Byte granular
 - Read
 - Extremely low latency
 - committed writes integrated on index traversal
 - Write
 - Incoming data and metadata logged
 - Integration processes inserts into index



Version Intent Logs





Summary

- Ubiquitous NVRAM changes the game
- 3 order of magnitude step change in performance from disk
 - Terabytes/s -> Petabytes/s
 - mS latency -> μS latency
- Workflows will change to exploit
 - Persistent Memory programming models
 - Data aware workflow scheduling
- Storage software must change to exploit
 - Same transactional guarantees required
 - End-to-end OS bypass required
 - Scalable comms/security context establishment
 - More I/O stack configuration flexibility





