



Advancing Digital Storage Innovation



Improvements in Lustre Data Integrity

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- Lustre Wire Checksum Improvements
 - Cleanup
 - Portability
 - Algorithms
 - Performance
- End-to-End Data Integrity
 - T10DIF/DIX
 - Version Mirroring

Goals

- End user assurance that their data was written to disk accurately
- Protection against all RAID (single) failure modes
- Offload the heavy calculations from the servers
- Support a wider range of client/server hardware

Over-the-Wire Bulk Checksums

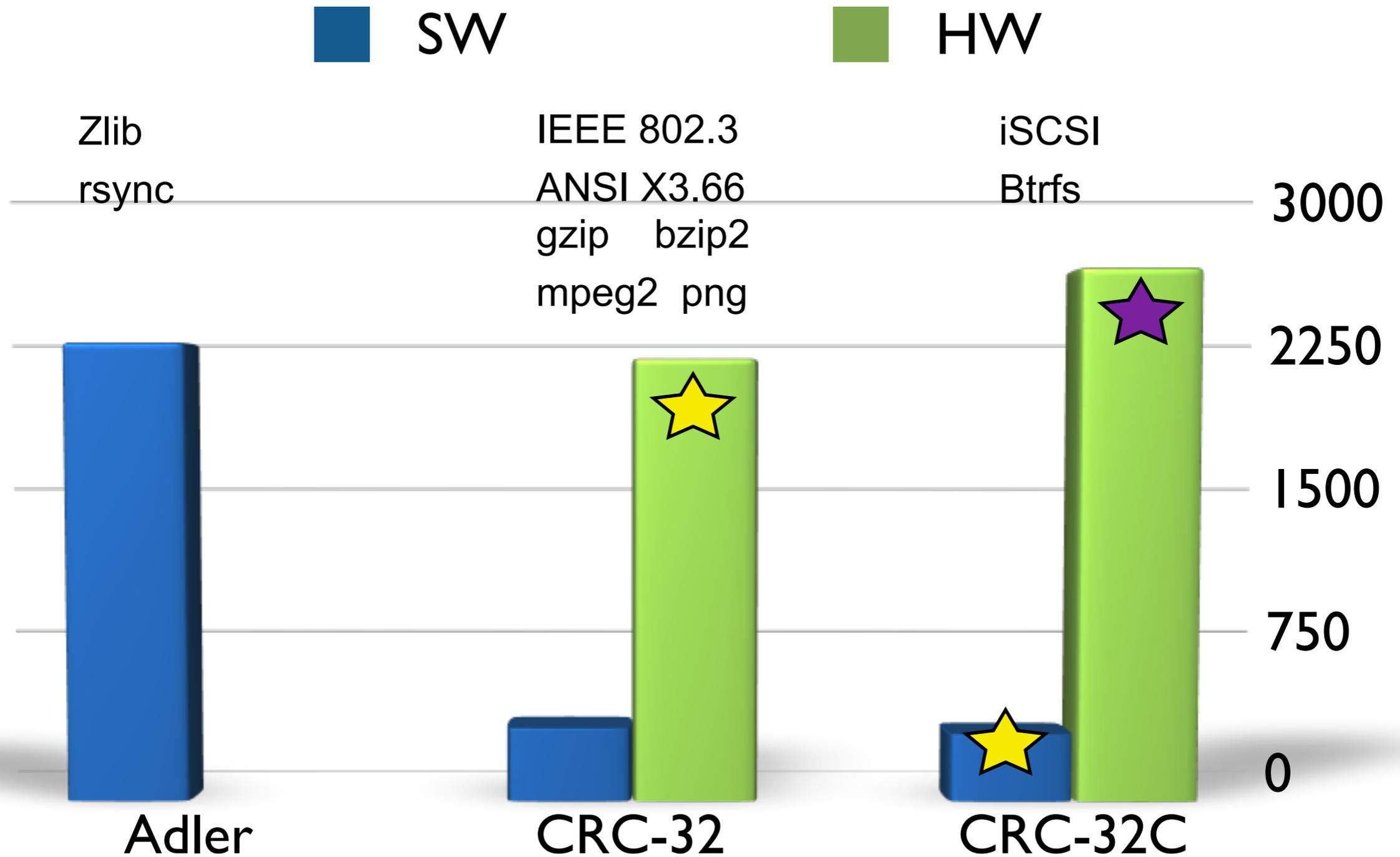
Bulk Checksum History

- Initially only software CRC-32 (IEEE, ANSI) 2007
- Adler-32 added in 1.6.5
 - easy to calculate
 - weak for small message sizes
- Shuichi Ihara added initial support for hardware CRC-32C (Castagnoli)
 - Intel Nehalem
 - bz 23549, landed in Lustre 2.2
- WC added multi-threaded ptlrpcd, and bulk RPC checksums moved into ptlrpcd context: parallelized checksums
 - LU-884, LU-1019, in Lustre 2.2
- sptlrpc implementation used a different set of functions
 - CRC-32, Adler, MD5, SHA1-512

Bulk Checksum Changes

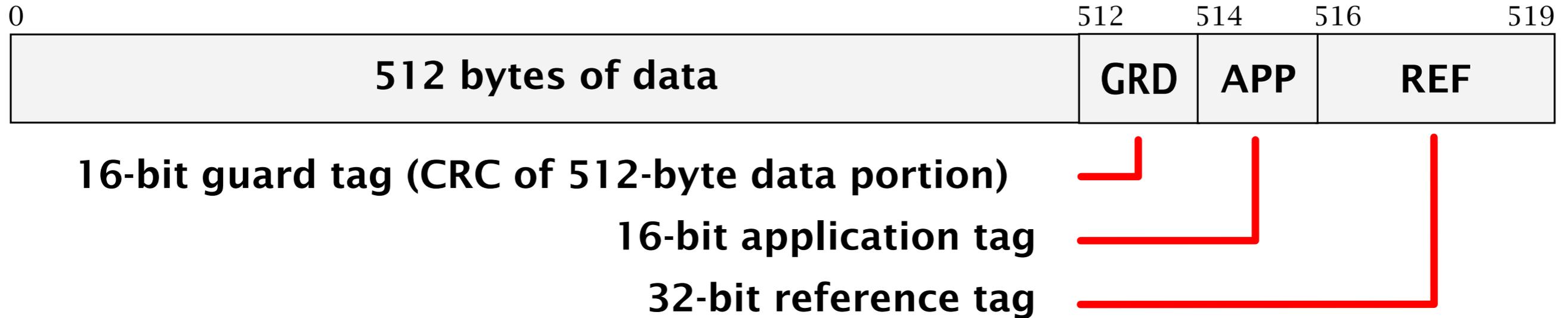
- Cleanup of sptlrpc and bulk checksum algorithms to use the kernel crypto hash library
 - simplifies future additions
 - LU-1201
- Addition of Software CRC-32C support
 - eg. if server has HW support and clients don't
 - LU-1201
- Implementation of Hardware CRC-32 using PCLMULQDQ
 - Intel Westmere
 - MRP-314, still testing

Bulk Checksum Speeds, MB/s



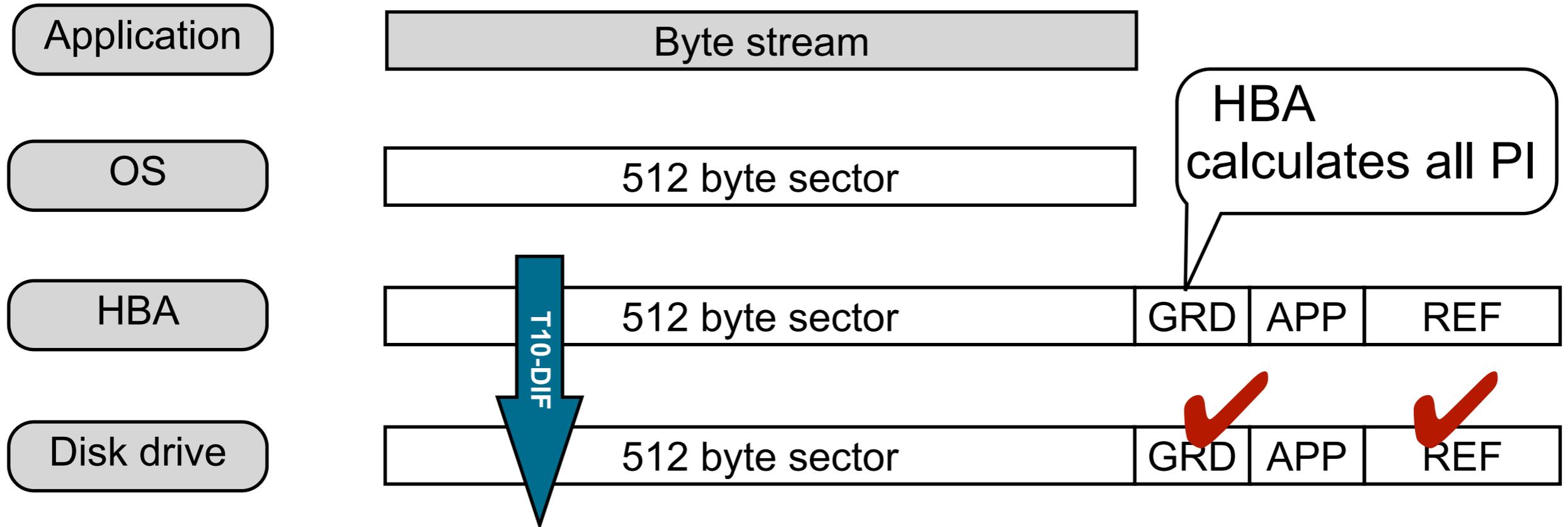
End-to-End Data Integrity with T10 and Version Mirroring

T10 DIF

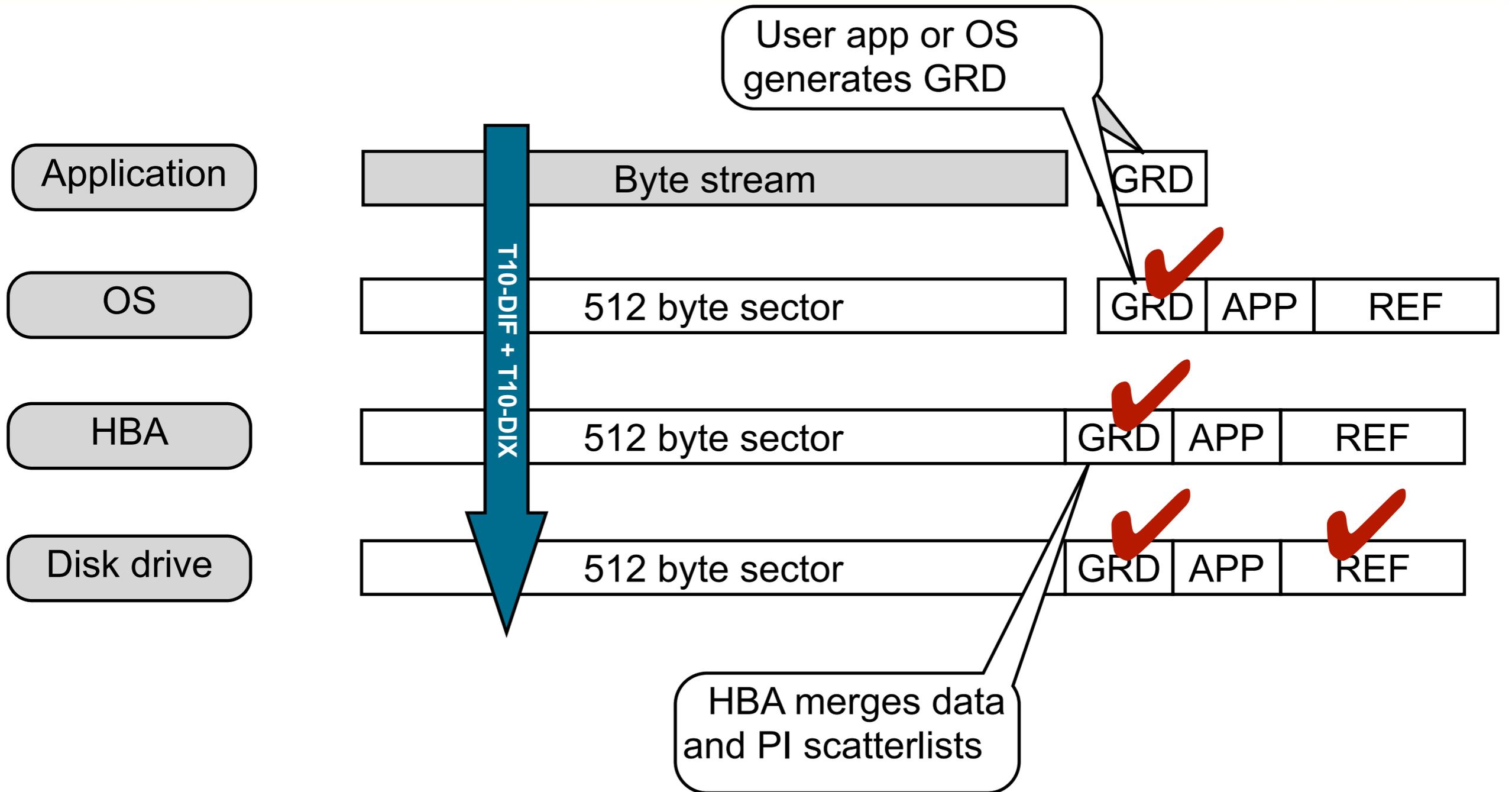


- 8 bytes of PI appended to 512 byte sectors
- HBA and disk drives must support T10-DIF in hardware

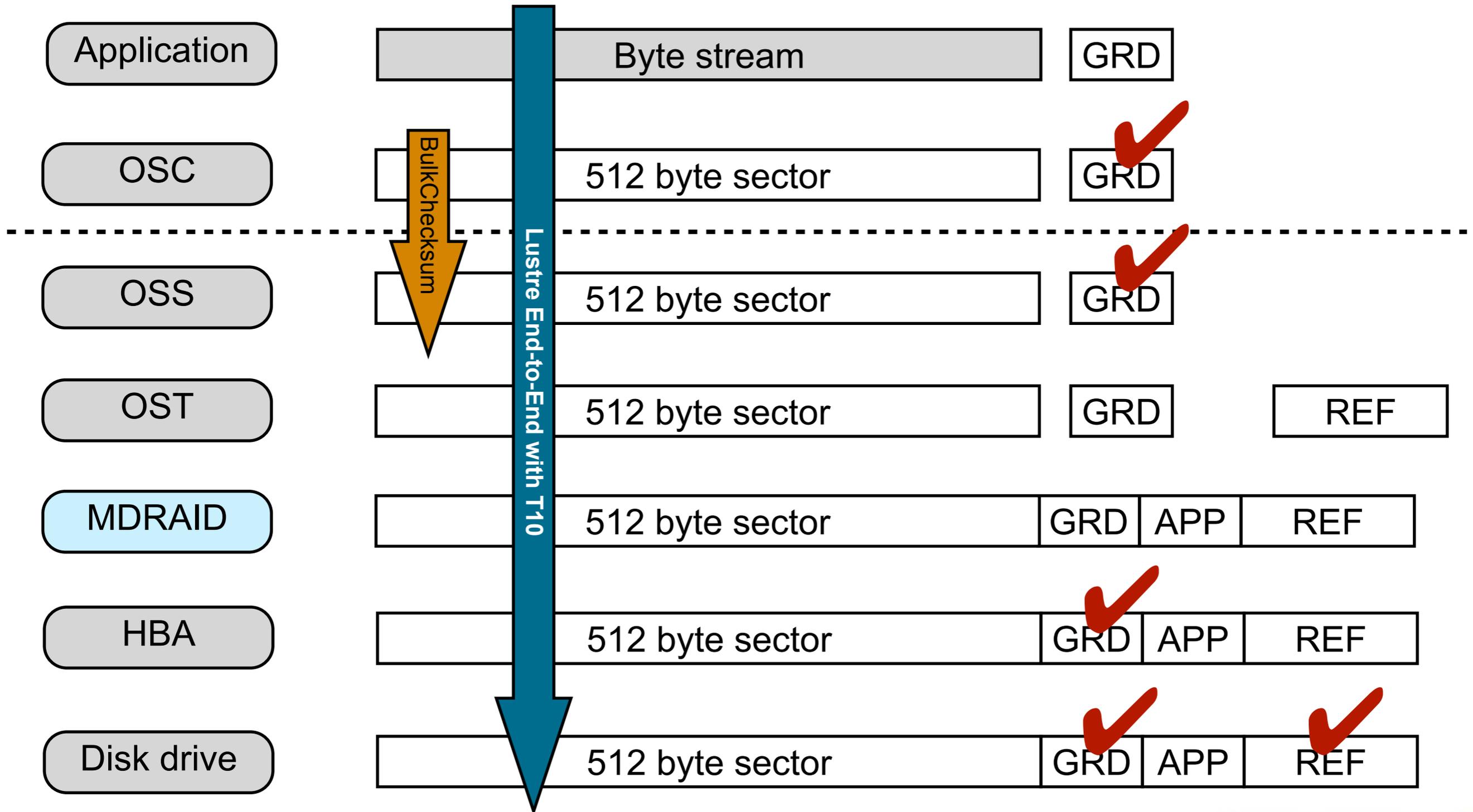
T10 DIF



T10 DIX



T10 DIX with Lustre



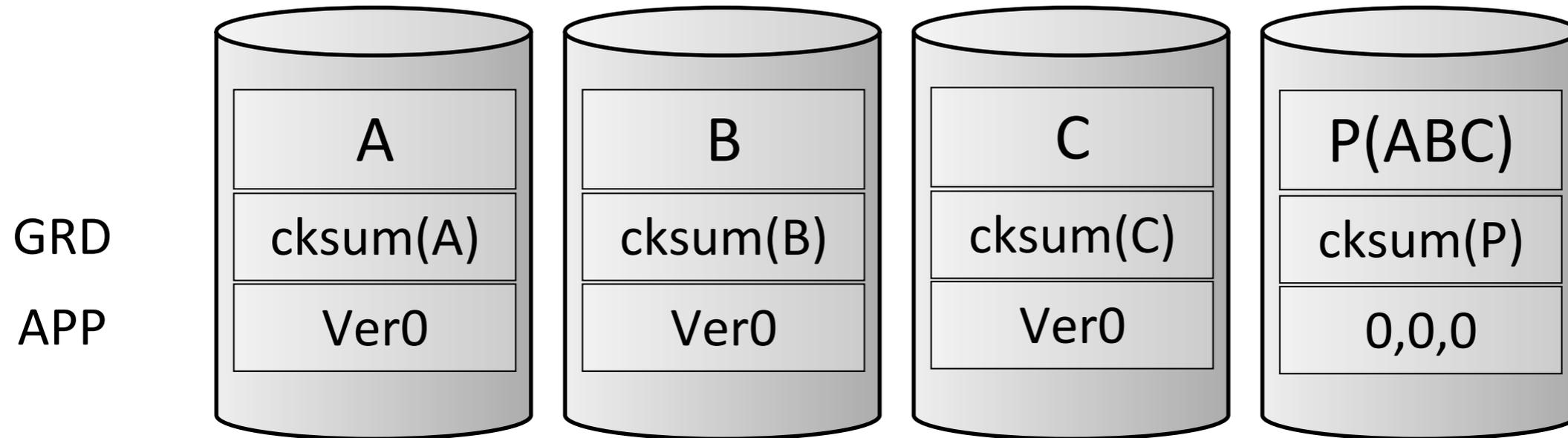
Changes to Lustre

- Additional checksum data described or carried in brw RPC
- Add PI and checking to data path
- For mmap'ed pages, early GRD failure implies data has changed, recompute from OSC
- Disable bulk checksums
- Optional GRD checking on OSS can push all checksum load to HBA/disk hardware

RAID failure modes

- ◉ *Parity Lost and Parity Regained* - Andrew Krioukov
 - Latent Sector (reliable read) errors
 - Data Corruption
 - Misdirected Writes
 - Lost Writes
 - Torn Writes
 - Parity Pollution
- Outcomes
 - Data recovery
 - Data loss (detected)
 - Data corruption (silent)

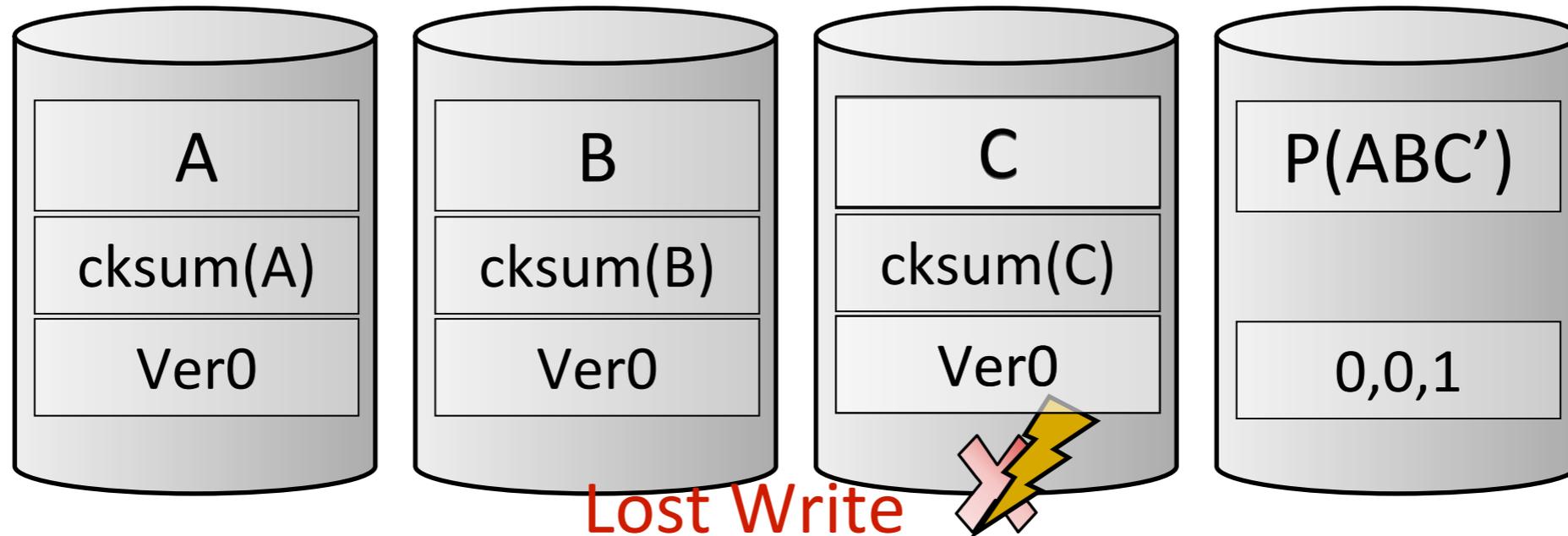
RAID with Version Mirroring



- RAID stripe across disks
- Block (chunk) on one disk
- Multiple sectors per block
- Store block version in T10 APP field
 - sectors within a chunk store the same version
 - parity block contains version vector

Rebuild with Version Mirroring

Update C to C', write C' and P(ABC')



Later, attempt to update A to A'

First, read B & C to prepare for constructing new P

But first read P to verify versions before writing P(A'BC')

Version mismatch, latest is in P, so reconstruct C' from P

$$C' = P(ABC') \oplus A \oplus B$$

Write C'

Calculate P(A'BC')

Write A' and P(A'BC')

RAID failure modes

- Latent Sector (reliable read) errors - drive detects
- Data Corruption - GRD, lightweight, partial reads
- Misdirected Writes - REF
- Lost Writes - parity block version vector
- Torn Writes - sector versions
- Parity Pollution - GRD + versions allow safe reconstruction

Fin