HPCS I/O Scenarios

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"Comparing system performance tools that measure I/O with I/O benchmarks is like comparing apples with flying pigs."

Henry Newman personal communication April 23, 2012





What are the HPCS Scenarios?

- As part of its High Productivity Computing Systems (HPCS) program, DARPA's HPCS Mission Partners provided vendors with a set of 14 *scalable* application workloads
 "HPCS Mission Partner File I/O Scenarios"
 <u>http://sourceforge.net/projects/hpcs-io/files/</u> DARPA.HPCS.IO.Scenarios.2011.pdf
- The Scenarios are not by themselves benchmarks; instead, they provide well-defined rules for objectively evaluating storage system scalability
- For all Scenarios, scaling performance is more important than absolute throughput for a single configuration
- The Scenarios are meant to show how a storage solution will scale when deployed with real workloads





HPCS I/O Environments



HPCS Mission Partners preferred one of these two I/O Environments



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Capture Environment

Parallel Environment

HPCS I/O Scenarios

- **1.** Single stream with large data blocks operating in half duplex mode
- 2. Single stream with large data blocks operating in full duplex mode
- 3. Multiple streams with large data blocks operating in full duplex mode
- 4. Extreme file creation rates
- 5. Checkpoint/restart with large I/O requests
- 6. Checkpoint/restart with small I/O requests
- 7. Checkpoint/restart Large file count per directory large I/Os
- 8. Checkpoint/restart large file count per directory small I/Os
- 9. Walking through directory trees
- **10.** Parallel walking through directory trees
- **11.** Random stat() system call to files in the file system (one process)
- **12.** Random stat() system call to files in the file system (multiple proc's)
- **13.** Small block random I/O to multiple files
- 14. Small block random I/O to a single file
- Lustre operates best in Parallel I/O Environment
- But don't forget there are customers for Capture I/O



Parallel file system workloads

- Characterizing parallel file systems can generally be summarized by evaluating the impact of the following on scalable I/O performance:
 - File access method
 - shared (N:1)
 - file per process (N:N)
 - I/O size
 - large [1 MB and 10 MB] (eg, segmented shared files)
 - small [32 KB and 128 KB](eg, strided shared files)
 - File access pattern
 - sequential
 - random
 - Metadata IOPs





HPCS Checkpoint/Restart Scenarios

Scenarios 5-8 vary by I/O record size and file access pattern

Scenario	I/O Record Size	File Access Pattern	
5	Large 1% of node memory		
6	Small 10,000 * (Real*8)	Shared (N-1) Strided access	
7	Large 1% of node memory	File per Process (N-N)	
8	Small 10,000 * (Real*8)	File per Process (N-N)	

 Each Scenario runs on at least three (3) different hardware configurations to demonstrate system scalability using hardware bottleneck analysis





Shared File Scenarios 5 & 6



Scenario 5 : Shared, Segmented I/O

- M processes (PEs) share one file
- File organized by PE
- Each PE creates 50 arrays
- Each array has N Real*8 elements
- All 50 arrays on each PE use 50% of node memory
- Each PE seeks to its location in the file based on its PE index
- Each PE writes each array of *N* elements sequentially from this offset

Scenario 6 : Shared, Strided I/O

- *M* processes (PEs) share one file
- · File organized by Array
- Each PE creates 50 arrays
- Each array has Z blocks of 10,000 Real*8 elements
- All 50 arrays on each PE use 25% of node memory
- For each array, each PE seeks a new location in the file based on the Array index
- Each PE writes the *Z* blocks of 10,000 elements sequentially from this offset





HPCS I/O Scenarios Tests

- DARPA provided the Scenarios, not the actual tests
 - A number of proprietary implementations of the Scenarios exist, but none have been available as open source
- Cray implemented the Scenarios for its HPCS program and published the source on SourceForge under Cray's BSDcompliant Open Source License
 - http://hpcs-io.cray.com/
 - The Scenarios and Cray's implementation are file system agnostic
 - The repository includes scripts with example command line parameters for running each test





Testing on ORNL's Cray XT4

In March 2011, after decommissioning their Cray XT4, Oak Ridge National Laboratory (ORNL) provided Cray dedicated access to the supercomputer and its file system to validate our initial implementation of the HPCS Scenarios tests



The system had

- 18 DDN 9550 storage controllers
- 72 Lustre object storage servers (OSS), running on XT4 IO nodes





Scalable Storage Unit (SSU) Definition



- Each SSU had 8 OSS nodes, 2 DDN 9550 racks, 16 OSTs
- Each OSS had two OSTs, one on each DDN rack
- Theoretical performance per SSU was ~5 GB/s





File System Definitions

 Cray reconfigured the storage to create three file systems, each with its own metadata server, using one, two and four SSUs

FS name	# of SSUs	# of OSSs	# of Racks	# of LUNs
FS1	1	8	2	16
FS2	2	16	4	32
FS4	4	32	8	64

- If the SSU was a bottleneck in a Scenario test, then repeating the test on a file system with twice the number of SSUs should double the performance of the Scenario
- The XT4 had enough Lustre client nodes and sufficient network bandwidth that we could run tests against all three file systems simultaneously





Client scalability of the three file systems



- Used Scenario 7 (N-N, sequential large I/Os) to find minimum client count to exercise each file system
- 4 clients per OST gave best performance





Large I/O with Scenarios 5 (N-1) and 7 (N-N)



- N-N (scenario 7) performance was greater than N-1 (Scenario 5: shared file, segmented access)
- For both of these sequential workloads performance increases with file system size, though N-1 does not scale as well as N-N





Random I/O with Scenario 13 (N-N)



- All three file systems showed extremely poor performance (<2 MiB/s) with the random I/O workloads
- There is no indication of performance scaling with increasing file system size





Metadata performance with Scenario 9



- Lustre currently has a single metadata server, regardless of the size of the number of OSSs in the file system
- As a result, metadata performance shows no scaling as the capacity of the file systems increased





Lessons Learned at ORNL

- Reducing test execution time
 - Test completions based on memory transfer size are indeterminate
 - Most tests reached a steady-state bandwidth or IOP rate long before the test completed the specified data transfer
 - As a result, Cray restructured its code to exit the test and report results after a user-specified time

Creating a cold metadata cache

- Scenarios 9-12 are to be run with both a cold and warm cache
- Restarting the storage system to create a cold cache is too intrusive and too time consuming
- Cray found that creating a dirty cache by running the metadata tests against a different copy of the directory tree was as effective as and gave similar results
- DARPA updated the Scenarios document with these notes and cleaned up confusing sections of the Scenario descriptions based on Cray's implementation





Conclusions

- The important idea behind the HPCS Scenarios is the definition of scalable workloads that create a level playing field for comparing the capabilities and scalability of different storage systems
- Cray has demonstrated the scalability of its implementation and the effectiveness of the tests to expose the strengths and weaknesses of a storage subsystem
- Cray's Scenarios tests are not just for Lustre, but can be used to evaluate any other scalable file system
- Download the Scenarios document and tests from SourceForge : <u>http://hpcs-io.cray.com/</u>





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